

KEY

Student Name and ID Number

MATH 3012 Quiz 3, April 19, 2007, WTT

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1. Write the general solution of the advancement operator equation:

$$(A+2)^3(A-3)^2 f = 0.$$

$$f(n) = c_1(-2)^n + c_2 n(-2)^n + c_3 n^2(-2)^n + c_4 3^n + c_5 n 3^n$$

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2. Find a particular solution to the advancement operator equation:

$$(A^2 + 2A - 15)f(n) = 27(4)^n.$$

Try $f(n) = c \cdot 4^n$

$$\left. \begin{aligned} c \cdot 4^{n+2} + 2c \cdot 4^{n+1} - 15c \cdot 4^n &= 27 \cdot 4^n \\ 16c \cdot 4^n + 8c \cdot 4^n - 15c \cdot 4^n &= 27 \cdot 4^n \\ 9c &= 27 \end{aligned} \right\} c = 3$$

$$\underline{\underline{f(n) = 3(4)^n}}$$

3. Find the unique solution to the advancement operator equation:

$$(A^2 + 2A - 15)f(n) = 27(4)^n \text{ with } f(0) = 18 \text{ and } f(1) = 1.$$

$$A^2 + 2A - 15 = (A+5)(A-3)$$

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general solution

$$f(n) = c_1(-5)^n + c_2(3)^n + 3 \cdot 4^n$$

$$c_1 + c_2 + 3 = 18$$

$$-5c_1 + 3c_2 + 12 = 1$$

$$\underline{c_1 + c_2 = 15}$$

$$\underline{-5c_1 + 3c_2 = -11}$$

$$\left. \begin{aligned} -8c_1 &= -56 \\ c_1 &= 7 \end{aligned} \right\} c_2 = 8$$

$$f_n = 7(-5)^n + 8(3^n) + 3(4)^n$$

4. Write the Inclusion-Exclusion formula for d_n , the number of derangements of $\{1, 2, \dots, n\}$:

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$$d_n = \sum_{i=0}^n (-1)^i \binom{n}{i} (n-i)!$$

5. Use the formula in the preceding question to find the value of d_4 .

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$$d_4 = \binom{4}{0} 4! - \binom{4}{1} 3! + \binom{4}{2} 2! - \binom{4}{3} 1! + \binom{4}{4} 0!$$

$$= 24 - 24 + 6 \cdot 2 - 4 \cdot 1 + 1$$

$$= 12 - 4 + 1$$

$$= 9$$

6. Verify your answer to the previous question by listing all the derangements on $\{1, 2, 3, 4\}$.

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- | | | |
|---------|---------|---------|
| 2 1 4 3 | 3 1 4 2 | 4 1 2 3 |
| 2 3 4 1 | 3 4 1 2 | 4 3 1 2 |
| 2 4 1 3 | 3 4 2 1 | 4 3 2 1 |

7. A data file digraph_data.txt has been read for a digraph whose vertex set is $[6]$. The weights on the directed edges are shown in the matrix below. In the space to the right, apply Dijkstra's algorithm to find the distance from vertex 1 to all other vertices in the graph. Also, for each x , find a shortest path from 1 to x .

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W	1	2	3	4	5	6
1	0	54	58	17	22	97
2	60	0	28	9	19	8
3	46	24	0	19	9	12
4	16	36	40	0	8	73
5	23	29	30	3	0	47
6	19	8	82	16	28	0

2	3	4	5	6
(1,2) 54	(1,3) 58	(1,4), 17	(1,5) 22	(1,6) 97
(1,4,2) 53	(1,4,3) 57		(1,5) 22	(1,4,6) 90
(1,5,2) 51	(1,5,3) 52			(1,5,6) 69
	(1,5,3) 52			(1,5,2,6) 59
				(1,5,2,6) 59

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8. The data file for a graph with vertex set $\{1, 2, \dots, 7\}$ is shown below. In the space to the right, list in order the edges that would be found in carrying out Kruskal's algorithm (avoid cycles) and Prim's algorithm (build tree). Vertex 1 is the root.

graph1.txt

7

2 7 24

6 1 35

2 1 38

7 6 39

1 4 45

6 4 47

3 1 53

7 3 54

4 7 56

5 2 58

4 5 60

Kruskal

2 7 24

6 1 35

2 1 38

1 4 45

3 1 53

5 2 58

Prim

6 1 35

2 1 38

2 7 24

1 4 45

3 1 53

5 2 58

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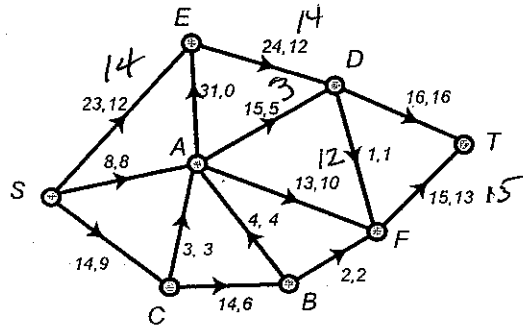
9. Let $R(n, m)$ denote the least positive integer t so that every graph on t vertices contains a complete subgraph of size n or an independent set of size m . Bob has a new computer and prides himself on being a good programmer. One day, he boasts to Alice that with the assistance of his new computer, he has succeeded in verifying that $R(100, 150) \leq 2^{300}$. Alice is not impressed. Can you explain why?

We know $R(n, m) \leq \binom{n+m-2}{n-1}$ so

$$R(100, 150) \leq \binom{248}{99} < 2^{248} < < < 2^{300}$$

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10.



- a. What is the current value of the flow? $12 + 8 + 9 = 29$
- b. What is the capacity of the cut $V = \{S, A, C, E\} \cup \{B, D, F, T\}$. $24 + 15 + 13 + 14 = 66$
- c. Carry out the labeling algorithm, using the pseudo-alphabetic order on the vertices and list below the labels which will be given to the vertices.

S (*, +, ∞)
 C (S, +, 5)
 E (S, +, 11)
 B (C, +, 5)
 D (E, +, 11)
 A (D, -, 5)
 F (A, +, 3)
 T (F, +, 2)

- d. Use your work in part c to find an augmenting path and make the appropriate changes directly on the diagram.

Augmenting path (S, E, D, A, F, T)

- e. What is the value of the new flow?
 $29 + 2 = 31$

Carry out the labeling algorithm a second time on the updated flow. It should halt without the sink being labeled. Find a cut whose capacity is equal to the value of the flow.

S (*, +, ∞)
 C (S, +, 5)
 E (S, +, 9)
 B (C, +, 5)
 D (E, +, 9)
 A (D, -, 3)
 F (A, +, 1)

Labeled

T

capacity =
 $16 + 15 = 31$

Unlabeled