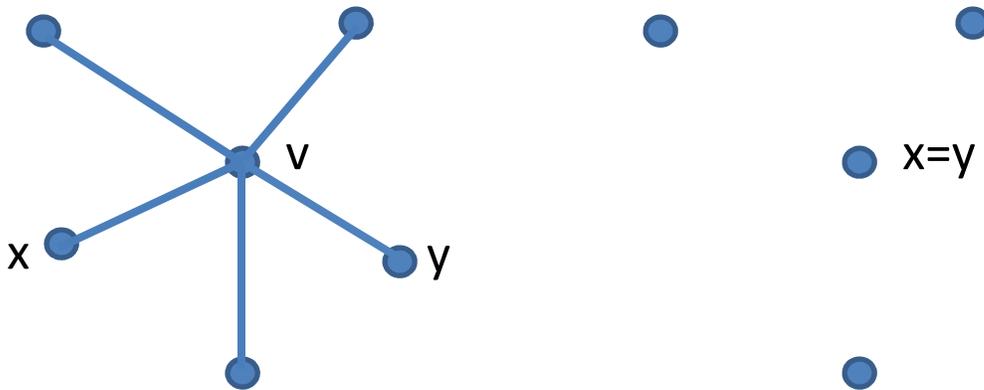


Corollary. Every simple planar graph G has a vertex of degree at most five.

Proof. We may assume G has $n \geq 3$ vertices. Then the sum of the degrees is $2|E(G)| \leq 6n - 12$ by Corollary 1.14, and hence G has a vertex of degree at most five.

Corollary. Every planar graph is 5-colorable.

Proof. Let v be a vertex of G of degree at most five. If v has degree at most four, then the result follows by induction applied to $G \setminus v$. So WMA v has degree five. Since G has no K_5 subgraph by planarity, some two neighbors of v , say x and y , are not adjacent. Apply induction to the graph obtained from $G \setminus v$ by identifying x and y .

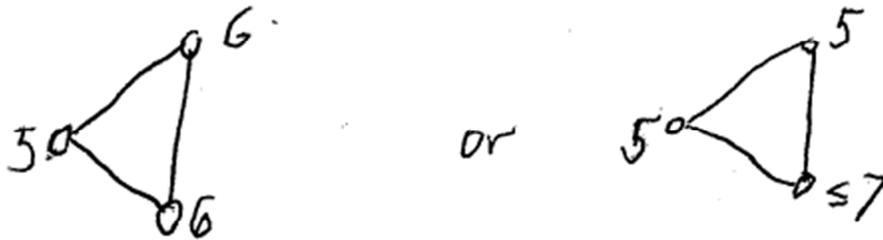


Theorem. Every planar graph is 4-colorable.

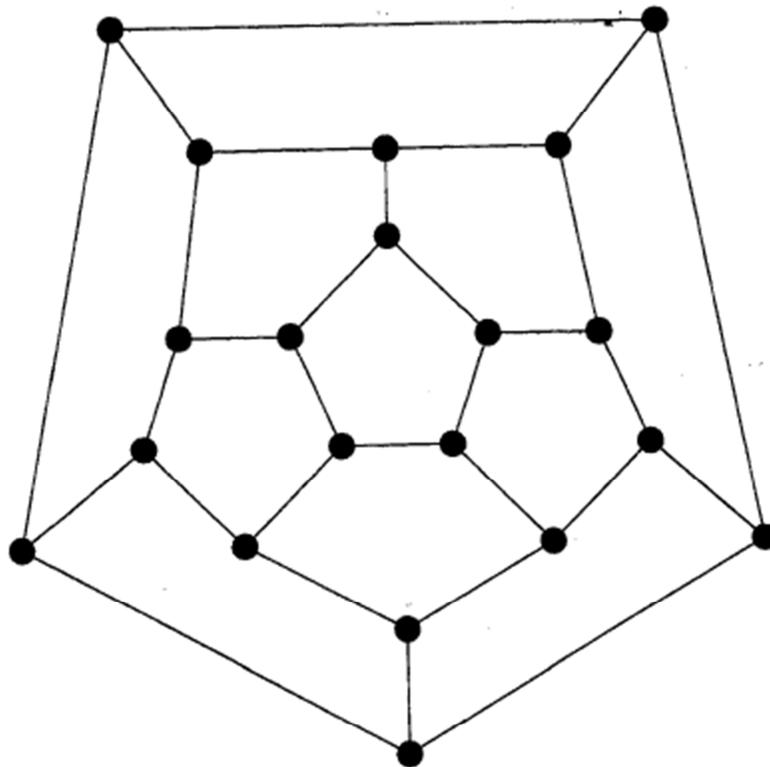
THE DISCHARGING METHOD

THEOREM. Every simple plane graph of minimum degree five has a face bounded by a triangle $C = xyz$ such that

$$\deg(x) + \deg(y) + \deg(z) \leq 17.$$

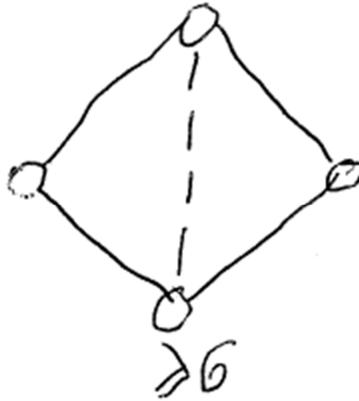


EXAMPLE. 17 is best possible



Dodecahedron

PROOF. We will allow G to be a multigraph, but insist every face have length ≥ 3 . Suppose for a contradiction that said face does not exist. We may assume that $E(G)$ is maximal. It follows that every vertex v incident with a ≥ 4 -face has degree five, for otherwise we can add an edge.



Assign charges as follows:

$$\text{ch}(v) := \deg(v) - 6$$

$$\text{ch}(f) := 2|f| - 6$$

Then the sum of all charges is

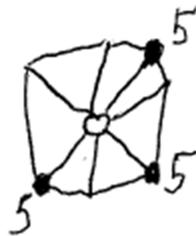
$$\begin{aligned} \sum \text{ch}(v) + \sum \text{ch}(f) &= \sum \deg(v) - 6|V(G)| + 2\sum |f| - 6|F(G)| = \\ &= 2|E(G)| - 6|V(G)| + 4|E(G)| - 6|F(G)| = \\ &= 6(|E(G)| - |V(G)| - |F(G)|) = -12 \end{aligned}$$

We now redistribute the charges as follows:

(1) Every ≥ 4 -face sends $1/2$ to every incident vertex (of degree five)



(2) Every 7-vertex sends $1/3$ to every neighbor of degree five



(3) Every ≥ 8 vertex v sends $1/4$ or 0 through every incident triangular face vxy :

if $\deg(x)=\deg(y)=5$, then both x,y receive $1/8$

if one of x,y has degree five, then it gets $1/4$

if x,y have degree ≥ 6 , then no charge is sent



Let $ch'(v)$ and $ch'(f)$ denote the new charges. Then

$$\sum ch'(v) + \sum ch'(f) = \sum ch(v) + \sum ch(f) = -12$$

But clearly $ch'(f) \geq 0$ for every face f .

We claim that $ch'(v) \geq 0$ for every vertex v , which gives a contradiction.

$ch'(v) = 0$ for every v of degree 6

$ch'(v) \geq 0$ for every v of degree 7, because every incident face is a triangle and at most 3 neighbors have degree 5.

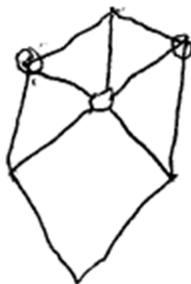
$ch'(v) \geq 0$ for every v of degree ≥ 8 , because it sends $\leq 1/4$ through every incident face

$ch'(v) \geq 0$ for every v of degree 5:

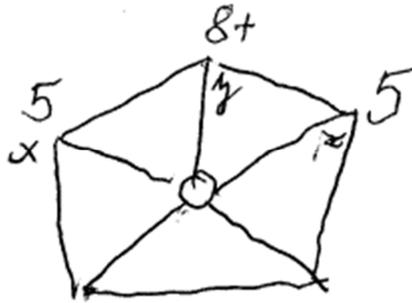
If at least two incident ≥ 4 -faces, then receive $1/2$ from each



If one incident ≥ 4 -face, then receive $1/2$ from it and at least two neighbors have degree ≥ 7 , each sending $\geq 1/4$



If every incident face is a triangle, then at least three neighbors have degree ≥ 7 . WMA one neighbor, say y , sends only $1/4$. Then x and z have degree 5, and the remaining two neighbors send $\geq 3/8$ each.



This proves the theorem.

Other uses:

$$\text{ch}(v) := 2\text{deg}(v) - 6$$

$$\text{ch}(f) := |f| - 6$$

Then the sum of the charges is -12

or

$$\text{ch}(v) := \text{deg}(v) - 4$$

$$\text{ch}(f) := |f| - 4$$

Then the sum of the charges is -8 .

Corollary. There is no planar graph of minimum degree at least 4 and every face of size at least 4.