# Reversal Ratio and Linear Extension Diameter 

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## Outline

(1) Linear Extension Diameter
(2) A Constant Bound?
(3) Posets of fixed dimension
(4) Posets of fixed width

## Linear Extension Diameter

## Definition

Let $\mathbf{P}$ be a finite poset. The linear extension graph $G(\mathbf{P})=(V, E)$ of $\mathbf{P}$ is defined as follows:

- $V$ is the set of all linear extensions of $\mathbf{P}$ and
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## Definition (Felsner and Reuter 1999)

The linear extension diameter of a finite poset $\mathbf{P}$, denoted $\operatorname{led}(\mathbf{P})$, is the diameter of its linear extension graph $G(\mathbf{P})$.

## Example



## Another Example



Felsner and Massow (2011)


## Reversal Ratio

## Definition

Let $\mathbf{P}$ be a poset and $L_{1}, L_{2}$ linear extensions of $\mathbf{P}$. We define the reversal ratio of the pair $\left(L_{1}, L_{2}\right)$ as

$$
R R\left(\mathbf{P} ; L_{1}, L_{2}\right)=\frac{\operatorname{dist}\left(L_{1}, L_{2}\right)}{\operatorname{inc}(\mathbf{P})} .
$$

The reversal ratio of $\mathbf{P}$ is

$$
R R(\mathbf{P})=\frac{\operatorname{led}(\mathbf{P})}{\operatorname{inc}(\mathbf{P})}=\max _{L_{1}, L_{2}} R R\left(\mathbf{P} ; L_{1}, L_{2}\right) .
$$

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- width $(\mathbf{P})=c|\mathbf{P}|$ for $c>0$ implies $R R(\mathbf{P})$ is large.


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- Unpublished example difficult to analyze.


## How small can $R R(\mathbf{P})$ be?

## Theorem (BK)

For every sufficiently large positive integer $k$, there exists a poset $\mathbf{P}_{k}$ of width $k$ with $R R\left(\mathbf{P}_{k}\right) \leq C / \log k$.

## Doubling Property

## Definition <br> Let $\mathbf{G}=(A \cup B, E)$ be a bipartite graph with $|A|=|B|=k$. We say that $\mathbf{G}$ has the doubling property if for every $Y \subset A$ with $|Y| \leq k / 3$, $|N(A)| \geq 2|A|$.

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## Lemma

Let $G_{d}(A, B)$ be a random $d$-regular bipartite graph on vertex sets $A$ and $B$ of size $k$, chosen according to the configuration model. For each $d \geq 10$ and $k$ sufficiently large, $G_{d}(A, B)$ has the doubling property with high probability.

## Construction of $\mathbf{P}_{k}$



## Analysis of $\mathbf{P}_{k}$

## Proposition

For $k$ sufficiently large and $\varepsilon \leq 1$ /logd, the number of incomparable pairs in $\mathbf{P}_{k}$ is at least

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\frac{r-2}{2(r-1)} \varepsilon^{2} k^{2} \log ^{2} k .
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## Proposition

For $k$ a sufficiently large multiple of three and $0<\varepsilon \leq 1$,

$$
\operatorname{led}\left(\mathbf{P}_{k}\right) \leq \frac{31}{6} \varepsilon k^{2} \log k .
$$

## Posets of fixed dimension

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For $d \geq 2$, define $\operatorname{DRR}(d)=\inf \{R R(\mathbf{P}): \operatorname{dim}(\mathbf{P})=d\}$.

## Bounding DRR(d)

- Standard example $\mathbf{S}_{n}$ ?


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$R R\left(\mathbf{n}^{d}\right) \rightarrow 1 / 2$ (Felsner-Massow)
- $\operatorname{DRR}(3)=2 / 3$ by considering $\mathbf{n}^{3}$
- $1 / 2 \leq \operatorname{DRR}(4) \leq 4 / 7$


## Dimension of $\mathbf{P}_{k}$

## Fact

The width of $\mathbf{P}_{k}$ is $k$, so $\operatorname{dim}\left(\mathbf{P}_{k}\right) \leq k$.

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## Corollary

For d sufficiently large,

$$
\operatorname{DRR}(d) \leq R R\left(\mathbf{P}_{d}\right) \leq \frac{27}{\log d} .
$$

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$W R R(3) \geq \operatorname{DRR}(3)$

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Proposition
WRR $(3) \leq 5 / 6$

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\begin{array}{cc}
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\hline L_{1} & L_{2} \\
\hline C_{4} & B_{4} / C_{4} \\
C_{3} / A_{4} / B_{4} & A_{4} \\
B_{3} & A_{3} / B_{3} / C_{3} \\
B_{2} / C_{2} / A_{3} & C_{2} \\
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- Bounds for $W R R(w)$ in general.


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## Conjecture (BK)

$$
W R R(3)=3 / 4
$$

## Thank You

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## Contact

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