

# category endomorphisms

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```
In[1]:= SetDirectory["1:"]; << goedel.09jul31a; << tools.m

:Package Title: goedel.09jul31a          2009 July 31 at 5:55 p.m.

It is now: 2009 Aug 6 at 9:58

Loading Simplification Rules

TOOLS.M                                Revised 2009 July 2

weightlimit = 40
```

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## summary

If  $u \cdot v$  is an endomorphism in a category, then  $v \cdot u$  is also an endomorphism.

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## endomorphisms

The class of **endomorphisms** in a category  $\text{cat}[x]$  is  $\text{image}[\text{hom}[\text{cat}[x]], \text{Id}] = \text{fix}[\text{domain}[\text{cat}[x]]]$ . For example, identity morphisms are endomorphisms:  $\text{ids}[\text{cat}[x]] \subset \text{fix}[\text{domain}[\text{cat}[x]]]$ . It will be shown below that one can equivalently characterize an endomorphism as a morphism that has equal **cod** and **dom**. The following variable-free statement of this is already available:

```
In[2]:= fix[composite[inverse[cod[cat[x]]], dom[cat[x]]]]

Out[2]= fix[domain[cat[x]]]
```

Theorem. The class of endomorphisms is contained in the class of all morphisms:

```
In[3]:= SubstTest[subclass, fix[t], domain[t], t -> domain[cat[x]]] // Reverse

Out[3]= subclass[fix[domain[cat[x]]], range[cat[x]]] == True

In[4]:= subclass[fix[domain[cat[x_]]], range[cat[x_]]] := True
```

Corollary. A simplification rule needed later.

```
In[5]:= equal[intersection[range[cat[x]], fix[domain[cat[x]]]], fix[domain[cat[x]]]]

Out[5]= True

In[6]:= intersection[fix[domain[cat[x_]]], range[cat[x_]]] := fix[domain[cat[x]]]
```

Theorem. An endomorphism has equal **cod** and **dom**.

```
In[7]:= Map[implies[#, equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], u]]] &,
  (member[u, fix[composite[inverse[funpart[y]], funpart[z]]]] // AssertTest) /.
  {y -> dom[cat[x]], z -> cod[cat[x]]}
```

```
Out[7]= or[equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], u]],
  not[member[u, fix[domain[cat[x]]]]] == True
```

```
In[8]:= or[equal[APPLY[cod[cat[x_]], u_], APPLY[dom[cat[x_]], u_]],
  not[member[u_, fix[domain[cat[x_]]]]] := True
```

Converse Theorem. A morphism with equal **cod** and **dom** is an endomorphism.

```
In[9]:= Map[implies[#, member[u, fix[domain[cat[x]]]]] &,
  (member[u, fix[composite[inverse[funpart[y]], funpart[z]]]] // AssertTest // Reverse) /.
  {y -> dom[cat[x]], z -> cod[cat[x]]}
```

```
Out[9]= or[member[u, fix[domain[cat[x]]]],
  not[equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], u]]],
  not[member[u, range[cat[x]]]] == True
```

```
In[10]:= or[member[u_, fix[domain[cat[x_]]]],
  not[equal[APPLY[cod[cat[x_]], u_], APPLY[dom[cat[x_]], u_]]],
  not[member[u_, range[cat[x_]]]] := True
```

---

## products of morphisms

If  $u \cdot v$  is the product of any two morphisms in a category  $x$ , then  $\text{pair}[u, v] \in \text{domain}[x]$ . If  $\text{pair}[u, v] \in \text{domain}[x]$  in a category  $x$ , then the **cod** of  $v$  is equal to the **dom** of  $u$ . Putting these two facts together yields the following general statement.

Theorem. If  $u \cdot v$  is any kind of morphism in a category  $x$ , then the **cod** of  $v$  is equal to the **dom** of  $u$ .

```
In[11]:= Map[not, SubstTest[and, implies[p1, p2], implies[p2, p3], not[implies[p1, p3]],
  {p1 -> member[APPLY[cat[x], PAIR[u, v]], y], p2 -> member[pair[u, v], domain[cat[x]]],
  p3 -> equal[APPLY[dom[cat[x]], u], APPLY[cod[cat[x]], v]]}] // Reverse
```

```
Out[11]= or[equal[APPLY[cod[cat[x]], v], APPLY[dom[cat[x]], u]],
  not[member[APPLY[cat[x], PAIR[u, v]], y]] == True
```

```
In[12]:= or[equal[APPLY[cod[cat[x_]], v_], APPLY[dom[cat[x_]], u_]],
  not[member[APPLY[cat[x_], PAIR[u_, v_]], y_]] := True
```

When the fact that endomorphisms have equal **cod** and **dom** is instantiated to the case of a product of two morphisms, rewrite rules about the **cod** and **dom** of a composite transform the statement as follows:

Lemma.

```
In[13]:= SubstTest[implies, member[t, fix[domain[cat[x]]]],
  equal[APPLY[cod[cat[x]], t], APPLY[dom[cat[x]], t]],
  t → APPLY[cat[x], PAIR[u, v]] // Reverse
```

```
Out[13]= or[equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], v]],
  not[member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]]],
  not[member[pair[u, v], domain[cat[x]]]]] = True
```

```
In[14]:= (% /. {u → u_, v → v_, x → x_}) /. Equal → SetDelayed
```

This result can be simplified. The composability literal is redundant.

Theorem. If  $u \cdot v$  is an endomorphism, then the **cod** of  $u$  is equal to the **dom** of  $v$ .

```
In[15]:= Map[not, SubstTest[and, implies[p1, p2], implies[and[p1, p2], p3], not[implies[p1, p3]],
  {p1 → member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]],
  p2 → member[pair[u, v], domain[cat[x]]],
  p3 → equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], v]]}] // Reverse
```

```
Out[15]= or[equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], v]],
  not[member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]]] = True
```

```
In[16]:= or[equal[APPLY[cod[cat[x_]], u_], APPLY[dom[cat[x_]], v_]],
  not[member[APPLY[cat[x_], PAIR[u_, v_]], fix[domain[cat[x_]]]]] := True
```

If the product of two morphisms is an endomorphism, then they can also be composed on the opposite order.

Corollary. If  $u \cdot v \in \text{fix}[\text{domain}[\text{cat}[x]]]$ , then  $v \cdot u$  is a morphism.

```
In[17]:= Map[not, SubstTest[and, implies[p1, p2],
  implies[p1, p3], implies[and[p2, p3], p4], not[implies[p1, p4]],
  {p1 → member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]],
  p2 → equal[APPLY[dom[cat[x]], v], APPLY[cod[cat[x]], u]],
  p3 → member[u, range[cat[x]]],
  p4 → member[pair[v, u], domain[cat[x]]]}] // Reverse
```

```
Out[17]= or[member[pair[v, u], domain[cat[x]]],
  not[member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]]] = True
```

```
In[18]:= or[member[pair[v_, u_], domain[cat[x_]]],
  not[member[APPLY[cat[x_], PAIR[u_, v_]], fix[domain[cat[x_]]]]] := True
```

Corollary. (Obtained by eliminating the variables  $u$  and  $v$ .)

```
In[19]:= Map[empty[domain[complement[#]]] &,
  complement[dif[image[inverse[cat[x]], fix[domain[cat[x]]]],
  inverse[domain[cat[x]]]]] // ReInNormality
```

```
Out[19]= subclass[inverse[image[inverse[cat[x]], fix[domain[cat[x]]]]], domain[cat[x]] = True
```

```
In[20]:= subclass[inverse[image[inverse[cat[x_]], fix[domain[cat[x_]]]]],
  domain[cat[x_]] := True
```

Corollary.

```
In[21]:= SubstTest[subclass, composite[Id, u], inverse[v],
  {u -> image[inverse[cat[x]], fix[domain[cat[x]]]], v -> domain[cat[x]]}] // Reverse
Out[21]= subclass[image[inverse[cat[x]], fix[domain[cat[x]]], inverse[domain[cat[x]]]] = True
In[22]:= subclass[image[inverse[cat[x_]], fix[domain[cat[x_]]]],
  inverse[domain[cat[x_]]]] := True
```

A counterexample will now be presented to show that this inclusion cannot be improved upon.

Lemma.

```
In[27]:= Map[not, SubstTest[implies, equal[u, v], equal[image[u, w], image[v, w]],
  {u -> SWAP, v -> composite[inverse[FIRST], SECOND], w -> id[set[0]]}] // Reverse
Out[27]= equal[SWAP, composite[inverse[FIRST], SECOND]] = False
In[28]:= equal[SWAP, composite[inverse[FIRST], SECOND]] := False
```

Counterexample. In general the inclusion derived above cannot be strengthened to an equation.

```
In[29]:= equal[image[inverse[cat[x]], fix[domain[cat[x]]], inverse[domain[cat[x]]]] /.
  x -> composite[SWAP, RIF]
Out[29]= False
```

The following is an instantiation to products for the converse statement that a morphism with equal **cod** and **dom** is an endomorphism.

Theorem.

```
In[30]:= SubstTest[or, member[t, fix[domain[cat[x]]]],
  not[equal[APPLY[cod[cat[x]], t], APPLY[dom[cat[x]], t]]],
  not[member[t, range[cat[x]]]], t -> APPLY[cat[x], PAIR[u, v]]] // Reverse
Out[30]= or[member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]],
  not[equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], v]]],
  not[member[pair[u, v], domain[cat[x]]]]] = True
In[31]:= or[member[APPLY[cat[x_], PAIR[u_, v_]], fix[domain[cat[x_]]]],
  not[equal[APPLY[cod[cat[x_]], u_], APPLY[dom[cat[x_]], v_]]],
  not[member[pair[u_, v_], domain[cat[x_]]]]] := True
```

The morphism variables **u** and **v** could be eliminated from this statement, but the same result can be obtained more simply as follows.

Lemma. An inclusion.

```
In[32]:= SubstTest[implies, subclass[u, v], subclass[image[t, u], image[t, v]],
  {t → cat[x], u → image[inverse[cat[x]], fix[domain[cat[x]]]},
  v → inverse[domain[cat[x]]]}] // Reverse
```

```
Out[32]= subclass[fix[domain[cat[x]]], image[cat[x], inverse[domain[cat[x]]]]] = True
```

```
In[33]:= (% /. x → x_) /. Equal → SetDelayed
```

It will be shown now that the reverse inclusion also holds.

Lemma.

```
In[34]:= Map[not, SubstTest[and, implies[p1, p2], implies[p1, p3],
  implies[and[p1, p2, p3], p4], implies[and[p1, p4], p5], not[implies[p1, p5]],
  {p1 → and[member[pair[u, v], domain[cat[x]]], member[pair[v, u], domain[cat[x]]]],
  p2 → or[member[pair[u, v], domain[cat[x]]], not[member[v, range[cat[x]]]],
  p3 → or[member[pair[u, v], domain[cat[x]]], not[member[u, range[cat[x]]]],
  p4 → or[equal[APPLY[cod[cat[x]], u], APPLY[dom[cat[x]], v]],
  not[member[pair[u, v], domain[cat[x]]]],
  p5 → member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]}] // Reverse
```

```
Out[34]= or[member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]],
  not[member[pair[u, v], domain[cat[x]]]],
  not[member[pair[v, u], domain[cat[x]]]]] = True
```

```
In[35]:= (% /. {u → u_, v → v_, x → x_}) /. Equal → SetDelayed
```

Eliminating the morphism variables  $u$  and  $v$  yields an inclusion:

Lemma.

```
In[36]:= Map[empty[domain[complement[#]]] &,
  SubstTest[class, pair[u, v], implies[member[pair[u, v], y], member[pair[u, v], z]],
  {y → intersection[domain[cat[x]], inverse[domain[cat[x]]]],
  z → image[inverse[cat[x]], fix[domain[cat[x]]]}] // Reverse
```

```
Out[36]= subclass[intersection[domain[cat[x]], inverse[domain[cat[x]]]],
  inverse[image[inverse[cat[x]], fix[domain[cat[x]]]]] = True
```

```
In[37]:= (% /. x → x_) /. Equal → SetDelayed
```

Corollary.

```
In[38]:= SubstTest[subclass, composite[Id, u], inverse[v],
  {u → intersection[domain[cat[x]], inverse[domain[cat[x]]]],
  v → image[inverse[cat[x]], fix[domain[cat[x]]]}] // Reverse
```

```
Out[38]= subclass[intersection[domain[cat[x]], inverse[domain[cat[x]]]],
  image[inverse[cat[x]], fix[domain[cat[x]]]] = True
```

```
In[39]:= (% /. x → x_) /. Equal → SetDelayed
```

An equation is obtained by combining this inclusion with the one in the opposite direction.

Theorem. An equation that can be made into a rewrite rule.

```
In[40]:= SubstTest[and, subclass[u, v], subclass[v, u],
  {u -> intersection[domain[cat[x]], inverse[domain[cat[x]]]],
  v -> image[inverse[cat[x]], fix[domain[cat[x]]]]}]

Out[40]= equal[image[inverse[cat[x]], fix[domain[cat[x]]]],
  intersection[domain[cat[x]], inverse[domain[cat[x]]]]] == True

In[41]:= image[inverse[cat[x_]], fix[domain[cat[x_]]]] :=
  intersection[domain[cat[x]], inverse[domain[cat[x]]]]
```

Corollary. A simpler equation.

```
In[42]:= ImageComp[cat[x], inverse[cat[x]], fix[domain[cat[x]]]] // Reverse

Out[42]= image[cat[x], inverse[domain[cat[x]]]] == fix[domain[cat[x]]]

In[43]:= image[cat[x_], inverse[domain[cat[x_]]]] := fix[domain[cat[x]]]
```

Corollary. If  $\mathbf{u} \cdot \mathbf{v}$  is an endomorphism in a category, then  $\mathbf{v} \cdot \mathbf{u}$  is also an endomorphism.

```
In[44]:= Map[not, SubstTest[and, implies[p1, p2], implies[p1, p3], implies[and[p2, p3], p4],
  not[implies[p1, p4]], {p1 -> member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]],
  p2 -> member[pair[v, u], domain[cat[x]]],
  p3 -> equal[APPLY[cod[cat[x]], v], APPLY[dom[cat[x]], u]],
  p4 -> member[APPLY[cat[x], PAIR[v, u]], fix[domain[cat[x]]]]}] // Reverse

Out[44]= or[member[APPLY[cat[x], PAIR[v, u]], fix[domain[cat[x]]]],
  not[member[APPLY[cat[x], PAIR[u, v]], fix[domain[cat[x]]]]] == True

In[45]:= or[member[APPLY[cat[x_], PAIR[v_, u_]], fix[domain[cat[x_]]]],
  not[member[APPLY[cat[x_], PAIR[u_, v_]], fix[domain[cat[x_]]]]] := True
```