

the class of connected relations

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```
In[1]:= SetDirectory["1:"]; << goedel.09oct29c; << tools.m

:Package Title: goedel.09oct29c          2009 October 29 at 3:55 p.m.

It is now: 2009 Oct 30 at 13:1

Loading Simplification Rules

TOOLS.M                                Revised 2009 October 28

weightlimit = 40
```

summary

The class **CONNEX** of connected relations is defined and its relation to the class of dichotomous relations is studied.

definition of CONNEX

A relation is **connected** if for every pair of distinct elements belonging to the union of its domain and range, one of them is related to the other. For example, if u and v are numbers, then the law of trichotomy implies that $u < v$ or $u = v$ or $u > v$. In this case, these three possibilities are mutually exclusive. For a connected relation in general the three possibilities need not be mutually exclusive.

The class of connected relations is defined by the following membership rule.

```
In[2]:= member[x_, CONNEX] :=
  and[member[x, V], subclass[x, cart[V, V]], subclass[cart[union[domain[x], range[x]],
    union[domain[x], range[x]]], union[Id, x, inverse[x]]]]
```

a formula for CONNEX

Theorem. A formula for the class of connected relations.

```
In[3]:= Map[empty, symdif[CONNEX, fix[
  composite[inverse[HULL[SYM]], S, IMAGE[id[Di]], CART, DUP, CUP, DORA]]] // Normality]

Out[3]= equal[CONNEX,
  fix[composite[inverse[HULL[SYM]], S, IMAGE[id[Di]], CART, DUP, CUP, DORA]]] == True
```

```
In[4]:= fix[composite[inverse[HULL[SYM]], S, IMAGE[id[Di]], CART, DUP, CUP, DORA] := CONNEX
```

There is a similar formula for the class of irreflexive connected relations.

Lemma.

```
In[5]:= (symdif[intersection[CONNEX, P[Di]],
          fix[composite[inverse[HULL[SYM]], IMAGE[id[Di]], CART, DUP, CUP, DORA]] //
          Normality) /. Equal -> SetDelayed
```

Theorem.

```
In[6]:= SubstTest[empty, symdif[u, v], {u -> intersection[CONNEX, P[Di]],
          v -> fix[composite[inverse[HULL[SYM]], IMAGE[id[Di]], CART, DUP, CUP, DORA]]}]
```

```
Out[6]= equal[fix[composite[inverse[HULL[SYM]], IMAGE[id[Di]], CART, DUP, CUP, DORA]],
           intersection[CONNEX, P[Di]]] = True
```

```
In[7]:= fix[composite[inverse[HULL[SYM]], IMAGE[id[Di]], CART, DUP, CUP, DORA] :=
          intersection[CONNEX, P[Di]]
```

dichotomous relations

A relation is **dichotomous** if for every pair of elements belonging to the union of its domain and range, one of them is related to the other. In the literature, such a relation is also called a **total relation** or a **strictly connected relation**. A dichotomous relation must be reflexive. When u and v are numbers, either $u \leq v$ or $v \leq u$ (or both). The two possibilities are not exclusive, but if both $u \leq v$ or $v \leq u$ hold for numbers, then $u = v$. In general however the definition of a dichotomous relation does not require this antisymmetric property. In the **GOEDEL** program, the class of dichotomous relations is called **DICHOT**. The membership rule for this class uses the fixed point class instead of the union of the domain and range.

```
In[8]:= member[x, DICHOT]
```

```
Out[8]= and[equal[cart[fix[x], fix[x]], union[x, inverse[x]]], member[x, V]]
```

Using **fix[x]** instead of **udora[x] = union[domain[x], range[x]]** makes no difference. To avoid needless duplication, the **GOEDEL** program has this rewrite rule:

```
In[9]:= equal[cart[udora[x], udora[x]], union[x, inverse[x]]]
```

```
Out[9]= equal[cart[fix[x], fix[x]], union[x, inverse[x]]]
```

Lemma. A dichotomous relation is connected.

```

In[10]:= Map[not, SubstTest[and, implies[p1, p2], implies[p2, p3],
  not[implies[p1, p3]], {p1 → equal[cart[fix[x], fix[x]], union[x, inverse[x]]],
  p2 → subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  union[x, inverse[x]]], p3 → subclass[cart[union[domain[x], range[x]],
  union[domain[x], range[x]]], union[Id, x, inverse[x]]]}] // Reverse

Out[10]= or[not[equal[cart[fix[x], fix[x]], union[x, inverse[x]]],
  subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  union[Id, x, inverse[x]]] == True

In[11]:= or[not[equal[cart[fix[x_], fix[x_]], union[inverse[x_], x_]],
  subclass[cart[union[domain[x_], range[x_]], union[domain[x_], range[x_]]],
  union[Id, inverse[x_], x_]] := True

```

Theorem.

```

In[12]:= Map[equal[V, #] &, SubstTest[class, x, not[member[x, t]], t → dif[DICHOT, CONNEX]]]

Out[12]= subclass[DICHOT, CONNEX] == True

In[13]:= subclass[DICHOT, CONNEX] := True

```

sum class

Lemma.

```

In[14]:= Map[empty, dif[CONNEX, P[cart[V, V]]] // Normality

Out[14]= subclass[U[CONNEX], cart[V, V]] == True

In[15]:= % /. Equal → SetDelayed

```

Lemma.

```

In[17]:= SubstTest[implies, subclass[u, v],
  subclass[U[u], U[v]], {u → DICHOT, v → CONNEX} // Reverse

Out[17]= subclass[cart[V, V], U[CONNEX]] == True

In[18]:= % /. Equal → SetDelayed

```

Theorem. Formula for the sum class of the class of connected relations.

```

In[20]:= SubstTest[and, subclass[u, v], subclass[v, u], {u → U[CONNEX], v → cart[V, V]}]

Out[20]= equal[cart[V, V], U[CONNEX]] == True

In[22]:= U[CONNEX] := cart[V, V]

```

reflexive connected relations

In this section it is shown that a relation is dichotomous if and only if it is reflexive and connected.

Lemma.

```
In[23]:= SubstTest[cliques, intersection[u, v],
             {u → union[Id, x, inverse[x]], v → cartsq[fix[x]]}]
Out[23]= intersection[chains[union[Id, x]], P[fix[x]]] == chains[x]
In[24]:= intersection[chains[union[Id, x_]], P[fix[x_]]] := chains[x]
```

Corollary.

```
In[25]:= SubstTest[subclass, P[fix[x]], intersection[u, v],
             {u → chains[union[Id, x]], v → P[fix[x]]}]
Out[25]= subclass[cart[fix[x], fix[x]], union[Id, x, inverse[x]]] ==
          subclass[cart[fix[x], fix[x]], union[x, inverse[x]]]
In[26]:= subclass[cart[fix[x_], fix[x_]], union[Id, x_, inverse[x_]]] :=
          subclass[cart[fix[x], fix[x]], union[x, inverse[x]]]
```

Because of a special rewrite rule for **fix[rfx[x]]**, a separate corollary is needed when the **rfx** wrapper is used.

Corollary. A temporary rewrite rule.

```
In[27]:= SubstTest[subclass, cart[fix[t], fix[t]],
             union[Id, t, inverse[t]], t → rfx[x] // Reverse
Out[27]= subclass[cart[fix[x], fix[x]], union[Id, inverse[rfx[x]], rfx[x]]] ==
          subclass[cart[fix[x], fix[x]], union[inverse[rfx[x]], rfx[x]]]
In[28]:= subclass[cart[fix[x_], fix[x_]], union[Id, inverse[rfx[x_]], rfx[x_]]] :=
          subclass[cart[fix[x], fix[x]], union[inverse[rfx[x]], rfx[x]]]
```

Eliminating the **rfx** wrapper yields the following lemma.

Lemma.

```
In[29]:= SubstTest[implies, and[equal[x, rfx[t]], subclass[cart[union[domain[x], range[x]],
             union[domain[x], range[x]], union[Id, x, inverse[x]]]],
             subclass[cart[fix[x], fix[x]], union[x, inverse[x]], t → x] // Reverse
Out[29]= or[not [REFLEXIVE[x]],
            not [subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]],
            union[Id, x, inverse[x]]]],
            subclass[cart[fix[x], fix[x]], union[x, inverse[x]]]]] == True
In[30]:= (% /. x → x_) /. Equal → SetDelayed
```

Theorem. A reflexive connected relation is dichotomous.

```
In[31]:= SubstTest[and, implies[p, subclass[u, v]], implies[p, subclass[v, u]],
  {p -> and[REFLEXIVE[x], subclass[cart[union[domain[x], range[x]],
    union[domain[x], range[x]]], union[Id, x, inverse[x]]]],
  u -> cart[fix[x], fix[x]], v -> union[x, inverse[x]]}]

Out[31]= or[equal[cart[fix[x], fix[x]], union[x, inverse[x]]], not[REFLEXIVE[x]],
  not[subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  union[Id, x, inverse[x]]]]] = True

In[32]:= or[equal[cart[fix[x_], fix[x_]], union[inverse[x_], x_]], not[REFLEXIVE[x_]],
  not[subclass[cart[union[domain[x_], range[x_]], union[domain[x_], range[x_]]],
  union[Id, inverse[x_], x_]]]] := True
```

Corollary. A concise variable-free restatement of the theorem (and its converse).

```
In[33]:= equal[intersection[RFX, CONNEX], DICHOT] // AssertTest

Out[33]= equal[DICHOT, intersection[CONNEX, RFX]] = True

In[34]:= intersection[CONNEX, RFX] := DICHOT
```

from dichotomous to connected and back

In this section it is shown how to obtain a connected relation from a dichotomous one and vice versa.

Lemma.

```
In[35]:= SubstTest[implies, equal[cartsq[udora[x]], y],
  subclass[cartsq[udora[intersection[Di, x]]], y], y -> union[x, inverse[x]] // Reverse

Out[35]= or[not[equal[cart[fix[x], fix[x]], union[x, inverse[x]]]],
  subclass[cart[union[fix[composite[Di, x]], fix[composite[x, Di]]],
  union[fix[composite[Di, x]], fix[composite[x, Di]]]], union[x, inverse[x]]]] = True

In[36]:= (% /. x -> x_) /. Equal -> SetDelayed
```

Theorem. If x is a dichotomous relation, then $\text{intersection}[Di, x]$ is connected.

```
In[37]:= Map[not, SubstTest[and, implies[p1, p2], implies[p2, p3],
  not[implies[p1, p3]], {p1 -> equal[cart[fix[x], fix[x]], union[x, inverse[x]]],
  p2 -> subclass[cart[union[fix[composite[Di, x]], fix[composite[x, Di]]],
  union[fix[composite[Di, x]], fix[composite[x, Di]]]], union[x, inverse[x]]],
  p3 -> subclass[cart[union[fix[composite[Di, x]], fix[composite[x, Di]]],
  union[fix[composite[Di, x]], fix[composite[x, Di]]]],
  union[Id, x, inverse[x]]]]] // Reverse

Out[37]= or[not[equal[cart[fix[x], fix[x]], union[x, inverse[x]]]],
  subclass[cart[union[fix[composite[Di, x]], fix[composite[x, Di]]], union[
  fix[composite[Di, x]], fix[composite[x, Di]]]], union[Id, x, inverse[x]]]] = True
```

```
In[39]:= or[not[equal[cart[fix[x_], fix[x_]], union[inverse[x_], x_]],
  subclass[cart[union[fix[composite[Di, x_]], fix[composite[x_, Di]]],
    union[fix[composite[Di, x_]], fix[composite[x_, Di]]]],
  union[Id, inverse[x_], x_]] := True
```

The following variable-free restatement is much easier to read.

Theorem. The strict part of a dichotomous relation is connected.

```
In[40]:= Map[equal[V, #] &, SubstTest[class, x, not[member[x, t]],
  t -> dif[DICHOT, image[inverse[IMAGE[id[Di]]], CONNEX]]]
```

```
Out[40]= subclass[image[IMAGE[id[Di]], DICHOT], CONNEX] == True
```

```
In[41]:= subclass[image[IMAGE[id[Di]], DICHOT], CONNEX] := True
```

One can go from a connected relation to a dichotomous relation by forming the reflexive closure.

Lemma. Simplification rule.

```
In[42]:= equal[image[HULL[RFX], DICHOT], DICHOT]
```

```
Out[42]= True
```

```
In[43]:= image[HULL[RFX], DICHOT] := DICHOT
```

Theorem. An inclusion in one direction.

```
In[44]:= SubstTest[implies, subclass[u, v], subclass[image[t, u], image[t, v]],
  {t -> HULL[RFX], u -> DICHOT, v -> CONNEX}] // Reverse
```

```
Out[44]= subclass[DICHOT, image[HULL[RFX], CONNEX]] == True
```

```
In[45]:= % /. Equal -> SetDelayed
```

To obtain an inclusion in the opposite direction, it is useful to temporarily introduce a variable.

Lemma.

```
In[46]:= SubstTest[implies, subclass[u, v],
  subclass[image[t, u], image[t, v]], {t -> id[cartsq[udora[x]]],
  u -> cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  v -> union[Id, x, inverse[x]]}] // Reverse
```

```
Out[46]= or[not[subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  union[Id, x, inverse[x]]],
  subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  union[composite[Id, x], id[union[domain[x], range[x]]], inverse[x]]]] == True
```

```
In[47]:= (% /. x -> x_) /. Equal -> SetDelayed
```

The expression **composite[Id, x]** in the above lemma can be viewed as a wrapper for a relation. Eliminating this wrapper yields the next lemma.

Lemma.

```
In[48]:= SubstTest[implies, equal[x, composite[Id, t]],
  implies[subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
    union[Id, x, inverse[x]]],
  subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
    union[x, id[union[domain[x], range[x]]], inverse[x]]], t -> x] // Reverse
```

```
Out[48]= or[not[subclass[x, cart[V, V]]],
  not[subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
    union[Id, x, inverse[x]]],
  subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
    union[x, id[union[domain[x], range[x]]], inverse[x]]] = True
```

```
In[49]:= (% /. x -> x_) /. Equal -> SetDelayed
```

The next lemma replaces an inclusion in the preceding lemma with an equation.

Lemma. The reflexive closure of a connected relation is dichotomous.

```
In[50]:= SubstTest[and, implies[p, subclass[u, v]],
  implies[p, subclass[v, u]], {p -> member[x, CONNEX],
  u -> cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  v -> union[x, id[union[domain[x], range[x]]], inverse[x]]}]
```

```
Out[50]= or[equal[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
  union[x, id[union[domain[x], range[x]]], inverse[x]]],
  not[member[x, V]], not[subclass[x, cart[V, V]]],
  not[subclass[cart[union[domain[x], range[x]], union[domain[x], range[x]]],
    union[Id, x, inverse[x]]]]] = True
```

```
In[51]:= (% /. x -> x_) /. Equal -> SetDelayed
```

Restatement.

```
In[52]:= implies[member[x, CONNEX], member[union[x, id[udora[x]]], DICHOT]]
```

```
Out[52]= True
```

To eliminate the variable x it is convenient to introduce a **setpart** wrapper.

Lemma. Introducing a **setpart** wrapper.

```
In[53]:= SubstTest[implies, member[t, CONNEX],
  member[union[t, id[udora[t]]], DICHOT], t → setpart[x]] // Reverse
```

```
Out[53]= or[equal[cart[union[domain[setpart[x]], range[setpart[x]]],
  union[domain[setpart[x]], range[setpart[x]]]],
  union[id[union[domain[setpart[x]], range[setpart[x]]]],
  inverse[setpart[x]], setpart[x]],
  not[subclass[cart[union[domain[setpart[x]], range[setpart[x]]],
  union[domain[setpart[x]], range[setpart[x]]]],
  union[Id, inverse[setpart[x]], setpart[x]]],
  not[subclass[setpart[x], cart[V, V]]] == True
```

```
In[54]:= (% /. x → x_) /. Equal → SetDelayed
```

Lemma. A simplification rule.

```
In[55]:= ImageComp[HULL[image[inverse[IMAGE[id[cart[V, V]]]], RFX]],
  id[P[cart[V, V]]], CONNEX] // Reverse
```

```
Out[55]= image[HULL[image[inverse[IMAGE[id[cart[V, V]]]], RFX], CONNEX] ==
  image[HULL[RFX], CONNEX]
```

```
In[56]:= % /. Equal → SetDelayed
```

Lemma. Elimination of the variable x .

```
In[57]:= Map[equal[V, #] &,
  SubstTest[class, x, implies[member[setpart[x], u], member[setpart[x], v]], {u → CONNEX,
  v → image[inverse[HULL[image[inverse[IMAGE[id[cart[V, V]]]], RFX]]], DICHOT}}]]
```

```
Out[57]= subclass[image[HULL[RFX], CONNEX], DICHOT] == True
```

```
In[58]:= % /. Equal → SetDelayed
```

The final step is the sharpen the inclusion to an equation.

Theorem. Every dichotomous relation is the reflexive closure of a trichotomous relation.

```
In[59]:= SubstTest[and, subclass[u, v], subclass[v, u],
  {u → image[HULL[RFX], CONNEX], v → DICHOT}]
```

```
Out[59]= equal[DICHOT, image[HULL[RFX], CONNEX]] == True
```

```
In[60]:= image[HULL[RFX], CONNEX] := DICHOT
```

Corollary.

```
In[61]:= SubstTest[implies, subclass[u, v], subclass[image[u, w], image[v, w]],
  {u → id[CONNEX], v → composite[inverse[HULL[RFX]], HULL[RFX]], w → CONNEX}] // Reverse
```

```
Out[61]= subclass[CONNEX, image[inverse[HULL[RFX]], DICHOT]] == True
```

```
In[62]:= subclass[CONNEX, image[inverse[HULL[RFX]], DICHOT]] := True
```