

# enumerating all ordinals in a given class: part 2. properties of partial enumerations

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```
In[1]:= SetDirectory["1:"]; << goedel.10oct04a

:Package Title: goedel.10oct04a          2010 October 4 at 9:30 a.m.

It is now: 2010 Oct 5 at 14:42

Loading Simplification Rules

TOOLS.M is now incorporated in the GOEDEL program as of 2010 September 3

weightlimit = 40
```

---

## summary

This is the second notebook in a series about enumerating all ordinals in a given class in increasing order. Here the focus is mainly on properties of partial enumerations.

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## monotonicity of partial enumerations

Theorem. Partial enumerations are strictly monotone.

```
In[2]:= Map[not, SubstTest[and, implies[p1, p2],
  implies[p2, p3], not[implies[p1, p3]], {p1 → member[w, partenum[x]],
  p2 → subclass[w, composite[HULL[intersection[x, OMEGA]], IMAGE[w]]],
  p3 → subclass[w, composite[S, IMAGE[w]]]}]] // Reverse

Out[2]= or[not[member[w, partenum[x]]], subclass[w, composite[S, IMAGE[w]]]] == True

In[3]:= or[not[member[w_, partenum[x_]]], subclass[w_, composite[S, IMAGE[w_]]]] := True
```

Corollary. Every partial enumeration is strictly monotone.

```
In[4]:= Map[not, SubstTest[and, implies[p1, p2], implies[p1, p3],
  implies[and[p2, p3], p4], not[implies[p1, p4]], {p1 → member[w, partenum[x]],
  p2 → FUNCTION[w], p3 → subclass[w, composite[S, IMAGE[w]]],
  p4 → subclass[composite[w, E, inverse[w]], E]}]] // Reverse

Out[4]= or[not[member[w, partenum[x]]], subclass[composite[w, E, inverse[w]], E]] == True

In[5]:= or[not[member[w_, partenum[x_]]], subclass[composite[w_, E, inverse[w_]], E]] := True
```

Lemma. A temporary simplification rule introduced to aid in the elimination of variables.

```
In[6]:= implies[member[w, partenum[x]], member[w, monotone[E, E]]] // NotNotTest
Out[6]= or[and[subclass[w, cart[V, V]], subclass[composite[w, E, inverse[w]], E]],
  not[member[w, partenum[x]]] == True

In[7]:= (% /. {w -> w_, x -> x_}) /. Equal -> SetDelayed
```

Theorem. Eliminating the variable  $w$ .

```
In[8]:= Map[equal[V, #] &, SubstTest[class, w,
  implies[member[w, u], member[w, v]], {u -> partenum[x], v -> monotone[E, E]}]]
Out[8]= subclass[partenum[x], monotone[E, E]] == True

In[9]:= subclass[partenum[x_], monotone[E, E]] := True
```

---

## ranges of partial enumerations

Lemma.

```
In[10]:= SubstTest[implies, subclass[x, t],
  subclass[range[x], range[t]], t -> composite[HULL[y], IMAGE[x]]] // Reverse
Out[10]= or[not[subclass[x, composite[HULL[y], IMAGE[x]]]],
  subclass[range[x], image[HULL[y], range[IMAGE[x]]]] == True

In[11]:= (% /. {x -> x_, y -> y_}) /. Equal -> SetDelayed
```

Theorem. The range of any partial enumeration is a set of ordinals.

```
In[12]:= Map[not, SubstTest[and, implies[p1, p2], implies[p2, p3],
  implies[p3, p4], not[implies[p1, p4]], {p1 -> member[w, partenum[x]],
  p2 -> subclass[w, composite[HULL[intersection[OMEGA, x]], IMAGE[w]]],
  p3 -> subclass[range[w], image[HULL[intersection[OMEGA, x]], range[IMAGE[w]]]],
  p4 -> subclass[range[w], OMEGA]}]] // Reverse
Out[12]= or[not[member[w, partenum[x]]], subclass[range[w], OMEGA]] == True

In[13]:= or[not[member[w_, partenum[x_]]], subclass[range[w_], OMEGA]] := True
```

Eliminating the variable  $w$  yields a similar result for  $\mathbf{enum}[x]$ .

Theorem. The range of  $\mathbf{enum}[x]$  is a class of ordinals (not necessarily a set).

```
In[14]:= Map[equal[V, #] &, SubstTest[class, w, implies[member[w, u], member[w, v]],
  {u -> partenum[x], v -> image[inverse[IMAGE[SECOND]], P[OMEGA]}]]]
Out[14]= subclass[range[enum[x]], OMEGA] == True
```

```
In[15]:= subclass[range[enum[x_]], OMEGA] := True
```

Corollary. The class `enum[x]` is a subclass of the cartesian square of  $\Omega$ .

```
In[16]:= SubstTest[subclass, composite[Id, t], cartsq[OMEGA], t -> enum[x]] // Reverse
```

```
Out[16]= subclass[enum[x], cart[OMEGA, OMEGA]] == True
```

```
In[17]:= subclass[enum[x_], cart[OMEGA, OMEGA]] := True
```

Any strictly monotone function whose domain and range are classes of ordinals is one-to-one.

Theorem. The class of partial enumerations is a class of bijections.

```
In[18]:= SubstTest[implies,
  and[subclass[u, v], subclass[v, w]], subclass[u, w], {u -> partenum[x],
  v -> intersection[FUNS, monotone[E, E], P[cartsq[OMEGA]]], w -> BIJ} // Reverse
```

```
Out[18]= subclass[partenum[x], BIJ] == True
```

```
In[19]:= subclass[partenum[x_], BIJ] := True
```

Corollary. Each partial enumeration is one-to-one.

```
In[20]:= SubstTest[implies, and[member[w, u], subclass[u, v]],
  member[w, v], {u -> partenum[x], v -> BIJ} // MapNotNot // Reverse
```

```
Out[20]= or[FUNCTION[inverse[w]], not[member[w, partenum[x]]]] == True
```

```
In[21]:= or[FUNCTION[inverse[w_]], not[member[w_, partenum[x_]]]] := True
```

## an alternate form of strict monotonicity

Lemma.

```
In[22]:= Map[subclass[#, monotone[complement[S], complement[S]]] &,
  SubstTest[intersection, P[cart[w, w]], monotone[restrict[t, w, w], restrict[t, w, w]],
  {t -> complement[S], w -> OMEGA}] // Reverse
```

```
Out[22]= subclass[intersection[monotone[E, E], P[cart[OMEGA, OMEGA]]],
  monotone[complement[S], complement[S]]] == True
```

```
In[23]:= subclass[intersection[monotone[E, E], P[cart[OMEGA, OMEGA]]],
  monotone[complement[S], complement[S]]] := True
```

Theorem. An alternate formulation of strict monotonicity for partial enumerations.

```
In[24]:= SubstTest[implies, and[subclass[u, v], subclass[v, w]], subclass[u, w],
  {u → partenum[x], v → intersection[monotone[E, E], P[cart[OMEGA, OMEGA]]],
  w → monotone[complement[S], complement[S]]} // Reverse
```

```
Out[24]= subclass[partenum[x], monotone[complement[S], complement[S]]] == True
```

```
In[25]:= subclass[partenum[x_], monotone[complement[S], complement[S]]] := True
```

Corollary.

```
In[27]:= SubstTest[implies, and[member[w, u], subclass[u, v]], member[w, v], {u → partenum[x],
  v → monotone[complement[S], complement[S]]} // Reverse // MapNotNot
```

```
Out[27]= or[not[member[w, partenum[x]]], subclass[composite[inverse[w], S, w], S]] == True
```

```
In[28]:= or[not[member[w_, partenum[x_]]], subclass[composite[inverse[w_], S, w_], S]] := True
```

---

## recursion equations

Because any partial enumeration is a function, the recursion inclusion can be sharpened by introducing a new variable to an equation involving function application. The following general result will be used.

Theorem. The value of a function at any point of its domain is the same as that of any extension.

```
In[29]:= SubstTest[implies, equal[x, t], equal[APPLY[x, z], APPLY[t, z]],
  t → composite[funpart[y], id[domain[x]]] // Reverse
```

```
Out[29]= or[equal[APPLY[x, z], APPLY[funpart[y], z]],
  not[member[z, domain[x]]], not[subclass[x, funpart[y]]] == True
```

```
In[30]:= or[equal[APPLY[x_, z_], APPLY[funpart[y_], z_]],
  not[member[z_, domain[x_]]], not[subclass[x_, funpart[y_]]] := True
```

The recursion condition  $w \subset \text{HULL}[x] \circ \text{IMAGE}[w]$  can be viewed as a statement that the function on the right side is an extension of  $w$ .

Lemma. Restatement of the recursion inclusion as an equation.

```
In[31]:= SubstTest[implies, and[member[t, domain[u]], subclass[u, funpart[v]],
  equal[APPLY[u, t], APPLY[funpart[v], t]], v → composite[HULL[x], IMAGE[u]]] // Reverse
```

```
Out[31]= or[equal[APPLY[u, t], hull[x, image[u, t]]], not[member[t, domain[u]]],
  not[subclass[u, composite[HULL[x], IMAGE[u]]]] == True
```

```
In[32]:= (% /. {t → t_, x → x_, u → u_}) /. Equal → SetDelayed
```

Theorem. The ordinal listed in at any point in a partial enumeration is the least ordinal in  $x$  greater than all ordinals listed previously.

```
In[33]:= Map[not, SubstTest[and, implies[p1, p2],
  implies[and[p0, p2], p3], not[implies[and[p0, p1], p3]],
  {p0 -> member[t, domain[w]], p1 -> member[w, partenum[x]],
  p2 -> subclass[w, composite[HULL[intersection[OMEGA, x]], IMAGE[w]]],
  p3 -> equal[APPLY[w, t], hull[intersection[OMEGA, x], image[w, t]]]}] // Reverse
```

```
Out[33]= or[equal[APPLY[w, t], hull[intersection[OMEGA, x], image[w, t]]],
  not[member[t, domain[w]]], not[member[w, partenum[x]]]] == True
```

```
In[34]:= or[equal[APPLY[w_, t_], hull[intersection[OMEGA, x_], image[w_, t_]]],
  not[member[t_, domain[w_]]], not[member[w_, partenum[x_]]]] := True
```

As a corollary, the following explicit formulation of the monotonicity condition for partial enumerations is obtained. Note that if  $t$  does not belong to the domain of  $w$ , then  $\text{APPLY}[w, t] = V$ , and the inclusion  $\text{image}[w, t] \subset \text{APPLY}[w, t]$  remains true.

Corollary. The ordinal listed at any point in a partial enumeration exceeds all previously listed ordinals.

```
In[35]:= Map[not, SubstTest[and, implies[and[p1, p2], p3], implies[p3, p4], or[p1, p4],
  not[implies[p2, p4]], {p1 -> member[t, domain[w]], p2 -> member[w, partenum[x]],
  p3 -> equal[APPLY[w, t], hull[intersection[OMEGA, x], image[w, t]]],
  p4 -> subclass[image[w, t], APPLY[w, t]]}] // Reverse
```

```
Out[35]= or[not[member[w, partenum[x]]], subclass[image[w, t], APPLY[w, t]]] == True
```

```
In[36]:= or[not[member[w_, partenum[x_]]], subclass[image[w_, t_], APPLY[w_, t_]]] := True
```

---

## truncating partial enumerations

In this section it is shown that the restriction of any partial enumeration to any ordinal is again a partial enumeration.

Lemma. A simplification rule.

```
In[37]:= SubstTest[subclass, composite[u, id[ord[v]]],
  composite[t, id[ord[v]]], t -> composite[x, IMAGE[id[ord[v]]]]]
```

```
Out[37]= subclass[composite[u, id[ord[v]]], composite[x, IMAGE[id[ord[v]]]]] ==
  subclass[composite[u, id[ord[v]]], x]
```

```
In[38]:= subclass[composite[u_, id[ord[v_]]], composite[x_, IMAGE[id[ord[v_]]]]] :=
  subclass[composite[u, id[ord[v]]], x]
```

Theorem. The restriction of any partial enumeration to any ordinal is a partial enumeration.

```

In[39]:= (Map[not, SubstTest[and, implies[and[p0, p2], p4],
  implies[and[p0, p3], p5], not[implies[and[p0, p1], p6]],
  {p0 -> equal[w, composite[u, id[ord[v]]], p1 -> member[u, partenum[x]],
  p2 -> subclass[u, composite[HULL[intersection[OMEGA, x]], IMAGE[u]]],
  p3 -> member[domain[u], OMEGA],
  p4 -> subclass[w, composite[HULL[intersection[OMEGA, x]], IMAGE[w]]],
  p5 -> member[domain[w], OMEGA], p6 -> member[w, partenum[x]]}] / .
  w -> composite[u, id[ord[v]]]) // Reverse

Out[39]= or[member[composite[u, id[ord[v]]], partenum[x]], not[member[u, partenum[x]]] == True

In[40]:= or[member[composite[u_, id[ord[v_]]], partenum[x_]],
  not[member[u_, partenum[x_]]] := True

```

Corollary. Removing the `ord` wrapper.

```

In[41]:= SubstTest[implies, equal[v, ord[t]], or[member[composite[u, id[v]], partenum[x]],
  not[member[u, partenum[x]]], t -> v] // Reverse

Out[41]= or[member[composite[u, id[v]], partenum[x]],
  not[member[u, partenum[x]]], not[member[v, OMEGA]] == True

In[42]:= or[member[composite[u_, id[v_]], partenum[x_]],
  not[member[u_, partenum[x_]]], not[member[v_, OMEGA]] := True

```

---

## domains of partial enumerations

The result of the preceding section is used to show that the class `image[IMAGE[FIRST], partenum[x]]` of domains of partial enumerations is a full subclass of  $\Omega$ .

Lemma.

```

In[43]:= SubstTest[implies, member[t, z],
  member[domain[t], image[IMAGE[FIRST], z]], t -> composite[x, id[y]] // Reverse

Out[43]= or[member[intersection[y, domain[x]], image[IMAGE[FIRST], z]],
  not[member[composite[x, id[y]], z]] == True

In[44]:= (% /. {x -> x_, y -> y_, z -> z_}) /. Equal -> SetDelayed

```

Theorem.

```

In[45]:= SubstTest[implies, equal[t, intersection[y, domain[x]]],
  or[member[t, image[IMAGE[FIRST], z]], not[member[composite[x, id[y]], z]],
  t -> y] // Reverse

Out[45]= or[member[y, image[IMAGE[FIRST], z]],
  not[member[composite[x, id[y]], z]], not[subclass[y, domain[x]]] == True

In[46]:= or[member[y_, image[IMAGE[FIRST], z_]],
  not[member[composite[x_, id[y_]], z_]], not[subclass[y_, domain[x_]]] := True

```

In the next lemma two steps are omitted to speed things up: `implies[and[p2, p3], p4]` and `implies[and[p5, p6], p7]`.

Lemma. Every ordinal less than the domain of some partial enumeration is also the domain of some partial enumeration.

```
In[47]:= Map[not,
  SubstTest[and, implies[p1, p3], implies[and[p1, p4], p5], implies[and[p2, p3], p6],
    not[implies[and[p1, p2], p7]], {p1 → member[u, partenum[x]],
      p2 → member[v, domain[u]], p3 → member[domain[u], OMEGA], p4 → member[v, OMEGA],
      p5 → member[composite[u, id[v]], partenum[x]], p6 → subclass[v, domain[u]],
      p7 → member[v, image[IMAGE[FIRST], partenum[x]]}] // Reverse
```

```
Out[47]= or[member[v, image[IMAGE[FIRST], partenum[x]]],
  not[member[u, partenum[x]], not[member[v, domain[u]]]] == True
```

```
In[48]:= (% /. {u → u_, v → v_, x → x_}) /. Equal → SetDelayed
```

Eliminating the variables `u` and `v` yields the following statement.

Lemma.

```
In[49]:= Map[empty[composite[Id, complement[#]]] &, SubstTest[class, pair[u, v],
  or[member[v, s], not[member[u, partenum[x]]], not[member[v, domain[u]]]],
  {s → image[IMAGE[FIRST], partenum[x]], t → partenum[x]}]
```

```
Out[49]= subclass[domain[enum[x]], image[IMAGE[FIRST], partenum[x]] == True
```

```
In[50]:= subclass[domain[enum[x_]], image[IMAGE[FIRST], partenum[x_]] := True
```

Restatement: The class of domains of partial enumerations is full.

```
In[51]:= full[image[IMAGE[FIRST], partenum[x]]]
```

```
Out[51]= True
```

Any full subclass of  $\Omega$  is either an ordinal or the class  $\Omega$  of all ordinals.

Corollary. The class of domains of partial enumerations for any class is either an ordinal or the class of all ordinals.

```
In[52]:= SubstTest[and, full[t], subclass[t, OMEGA], t → image[IMAGE[FIRST], partenum[x]]]
```

```
Out[52]= or[equal[OMEGA, image[IMAGE[FIRST], partenum[x]]],
  member[image[IMAGE[FIRST], partenum[x]], OMEGA]] == True
```

```
In[53]:= or[equal[OMEGA, image[IMAGE[FIRST], partenum[x_]]],
  member[image[IMAGE[FIRST], partenum[x_]], OMEGA]] := True
```