

monotonicity properties of exponentiation

Johan G. F. Belinfante
2007 January 18

```
In[1]:= SetDirectory["1:"]; << goedel89.13a; << tools.m

:Package Title: goedel89.13a      2007 January 13 at 9:15 p.m.

It is now: 2007 Jan 18 at 13:32

Loading Simplification Rules

TOOLS.M                          Revised 2007 January 7

weightlimit = 40
```

summary

In this notebook it is shown that for natural numbers, if $m \leq n$, then $m^p \leq n^p$. The converse holds if the power p is not zero. If the base b is not zero and if $m \leq n$, then $b^m \leq b^n$. The converse holds if $b > 1$. The proofs use only cardinality arguments and the laws of exponents, and do not require the use of induction. A key observation is that one can reduce inequalities to equations. Note that the inequality $m \leq n$ holds if and only if one can write n as m plus some natural number, and that number is in fact $n - m$.

```
In[2]:= equal[nat[x], natadd[nat[y], natsub[nat[x], nat[y]]]]

Out[2]= not[member[nat[x], nat[y]]]
```

fixed power

The monotonicity rule for fixed powers can be deduced most simply using a cardinality argument.

```
In[3]:= SubstTest[implies, subclass[fin[u], fin[v]], subclass[card[fin[u]], card[fin[v]]],
  {u → map[nat[z], nat[y]], v → map[nat[z], nat[x]]} // Reverse // MapNotNot

Out[3]= or[member[nat[x], nat[y]],
  not[member[natexp[nat[x], nat[z]], natexp[nat[y], nat[z]]]]] == True

In[4]:= or[member[nat[x_], nat[y_]],
  not[member[natexp[nat[x_], nat[z_]], natexp[nat[y_], nat[z_]]]]] := True
```

Conversely, one can show that if p is not zero and $m < n$, then $m^p < n^p$. The proof also does not require the use of induction. A lemma about mapping classes is needed:

```
In[5]:= SubstTest[and, subclass[u, v], subclass[v, u], {u → map[x, y], v → map[x, z]}]
Out[5]= equal[map[x, y], map[x, z]] == or[equal[0, x], equal[y, z], not[member[x, V]]]
In[6]:= equal[map[x_, y_], map[x_, z_]] := or[equal[0, x], equal[y, z], not[member[x, V]]]
```

Lemma.

```
In[7]:= SubstTest[implies, and[subclass[u, v], member[v, FINITE], equal[card[u], card[v]]],
  equal[u, v], {u → map[nat[z], nat[x]], v → map[nat[z], nat[y]]}] // Reverse // MapNotNot
Out[7]= or[equal[0, nat[z]], not[equal[natexp[nat[x], nat[z]], natexp[nat[y], nat[z]]]],
  not[member[nat[x], nat[y]]]] = True
In[8]:= (% /. {x → x_, y → y_, z → z_}) /. Equal → SetDelayed
```

The rule of trichotomy is also needed to obtain the converse theorem:

```
In[9]:= Map[not, SubstTest[and, implies[and[p1, p2], p3], implies[p1, p4],
  implies[p4, p5], implies[and[p3, p5], p6], not[implies[and[p1, p2], p6]],
  {p1 → member[nat[x], nat[y]], p2 → not[equal[0, nat[z]]],
  p3 → not[equal[natexp[nat[x], nat[z]], natexp[nat[y], nat[z]]]],
  p4 → not[member[nat[y], nat[x]]],
  p5 → not[member[natexp[nat[y], nat[z]], natexp[nat[x], nat[z]]]],
  p6 → member[natexp[nat[x], nat[z]], natexp[nat[y], nat[z]]]}] // Reverse
Out[9]= or[equal[0, nat[z]], member[natexp[nat[x], nat[z]], natexp[nat[y], nat[z]]],
  not[member[nat[x], nat[y]]]] = True
In[10]:= or[equal[0, nat[z_]], member[natexp[nat[x_], nat[z_]], natexp[nat[y_], nat[z_]]],
  not[member[nat[x_], nat[y_]]]] := True
```

fixed base $b = 2$

The special case that $m \leq n$ implies $2^m \leq 2^n$ can also be deduced using a cardinality argument.

```
In[11]:= SubstTest[implies, subclass[fin[u], fin[v]],
  subclass[card[fin[u]], card[fin[v]]], {u → P[nat[y]], v → P[nat[x]]}] // Reverse
Out[11]= or[member[nat[x], nat[y]],
  not[member[natexp[succ[set[0]], nat[x]], natexp[succ[set[0]], nat[y]]]]] = True
In[12]:= (% /. {x → x_, y → y_}) /. Equal → SetDelayed
```

Again, one can obtain a converse as follows.

```
In[13]:= SubstTest[implies, and[subclass[u, v], member[v, FINITE], equal[card[u], card[v]]],
  equal[u, v], {u → P[nat[x]], v → P[nat[y]]}] // Reverse
Out[13]= or[not[equal[natexp[succ[set[0]], nat[x]], natexp[succ[set[0]], nat[y]]]],
  not[member[nat[x], nat[y]]]] = True
```

```
In[14]:= (% /. {x → x_, y → y_}) /. Equal → SetDelayed
```

```
In[15]:= Map[not, SubstTest[and, implies[p1, p2], implies[p1, p3], implies[p3, p4],
  implies[and[p2, p4], p5], not[implies[p1, p5]], {p1 → member[nat[x], nat[y]],
  p2 → not[equal[natexp[succ[set[0]], nat[x]], natexp[succ[set[0]], nat[y]]]},
  p3 → not[member[nat[y], nat[x]]],
  p4 → not[member[natexp[succ[set[0]], nat[y]], natexp[succ[set[0]], nat[x]]], p5 →
  member[natexp[succ[set[0]], nat[x]], natexp[succ[set[0]], nat[y]]]}] // Reverse
```

```
Out[15]= or[member[natexp[succ[set[0]], nat[x]], natexp[succ[set[0]], nat[y]]],
  not[member[nat[x], nat[y]]] == True
```

```
In[16]:= (% /. {x → x_, y → y_}) /. Equal → SetDelayed
```

These two results can be combined into a single rewrite rule.

```
In[17]:= equiv[member[natexp[succ[set[0]], nat[x]], natexp[succ[set[0]], nat[y]]],
  member[nat[x], nat[y]]]
```

```
Out[17]= True
```

```
In[18]:= member[natexp[succ[set[0]], nat[x_]], natexp[succ[set[0]], nat[y_]]] :=
  member[nat[x], nat[y]]
```

general case of a fixed base

Theorem. If $m \leq n$, then $b^m \mid b^n$.

```
In[19]:= SubstTest[implies, equal[nat[x], natadd[nat[w], nat[y]]],
  member[pair[natexp[nat[z], nat[y]], natexp[nat[z], nat[x]]], DIV],
  w → natsub[nat[x], nat[y]] // Reverse
```

```
Out[19]= or[member[nat[x], nat[y]],
  member[pair[natexp[nat[z], nat[y]], natexp[nat[z], nat[x]]], DIV] == True
```

```
In[20]:= or[member[nat[x_], nat[y_]],
  member[pair[natexp[nat[z_], nat[y_]], natexp[nat[z_], nat[x_]]], DIV] := True
```

Lemma.

```
In[21]:= or[not[member[x, y]], not[member[pair[y, x], DIV]] // NotNotTest
```

```
Out[21]= or[not[member[x, y]], not[member[pair[y, x], DIV]] ==
  or[equal[0, y], not[equal[0, x]], not[member[y, omega]]]
```

```
In[22]:= or[not[member[x_, y_]], not[member[pair[y_, x_], DIV]] :=
  or[equal[0, y], not[equal[0, x]], not[member[y, omega]]]
```

Corollary. If b is not zero and if $m \leq n$, then $b^m \leq b^n$.

```
In[23]:= Map[not, SubstTest[and, implies[p2, p3],
  implies[and[p1, p3], p4], not[implies[and[p1, p2], p4]],
  {p1 → not[empty[nat[x]]], p2 → subclass[nat[z], nat[y]],
  p3 → member[pair[natexp[nat[x], nat[z]], natexp[nat[x], nat[y]]], DIV],
  p4 → subclass[natexp[nat[x], nat[z]], natexp[nat[x], nat[y]]]}] // Reverse
```

```
Out[23]= or[equal[0, nat[x]], member[nat[y], nat[z]],
  not[member[natexp[nat[x], nat[y]], natexp[nat[x], nat[z]]]] == True
```

```
In[24]:= or[equal[0, nat[x_]], member[nat[y_], nat[z_]],
  not[member[natexp[nat[x_], nat[y_]], natexp[nat[x_], nat[z_]]]] := True
```

Corollary.

```
In[25]:= SubstTest[or, equal[0, nat[w]], member[nat[x], nat[y]], not[
  member[natexp[nat[w], nat[x]], natexp[nat[w], nat[y]]], w → succ[nat[z]]] // Reverse
```

```
Out[25]= or[member[nat[x], nat[y]],
  not[member[natexp[succ[nat[z]], nat[x]], natexp[succ[nat[z]], nat[y]]]] == True
```

```
In[26]:= or[member[nat[x_], nat[y_]],
  not[member[natexp[succ[nat[z_]], nat[x_]], natexp[succ[nat[z_]], nat[y_]]]] := True
```

converse

For the converse, one needs to exclude the case $\mathbf{b} = 1$. It will be shown that if $\mathbf{b} > 1$ and $\mathbf{m} < \mathbf{n}$ then $\mathbf{b}^{\mathbf{m}} < \mathbf{b}^{\mathbf{n}}$. This proof again does not require the use of induction. This is known: if $\mathbf{b} > 0$ and $\mathbf{m} \leq \mathbf{n}$, then $\mathbf{b}^{\mathbf{m}} \leq \mathbf{b}^{\mathbf{n}}$. What we want is an implication going in the opposite direction. To begin, just replace \mathbf{x} with its successor. This yields:

```
In[27]:= SubstTest[implies, and[not[equal[0, nat[w]]], subclass[nat[t], nat[y]]],
  subclass[natexp[nat[w], nat[t]], natexp[nat[w], nat[y]]], t → succ[nat[x]]] // Reverse
```

```
Out[27]= or[equal[0, nat[w]], not[member[nat[x], nat[y]]],
  not[member[natexp[nat[w], nat[y]], natmul[nat[w], natexp[nat[w], nat[x]]]]] == True
```

```
In[28]:= (% /. {w → w_, x → x_, y → y_}) /. Equal → SetDelayed
```

Lemma.

```
In[29]:= equiv[and[member[set[0], nat[w]], not[equal[0, nat[w]]], member[set[0], nat[w]]]
```

```
Out[29]= True
```

```
In[30]:= and[member[set[0], nat[w_]], not[equal[0, nat[w_]]] := member[set[0], nat[w]]
```

Theorem. If $\mathbf{b} > 1$ then $\mathbf{b}^{\mathbf{m}} < \mathbf{b}^{\mathbf{m}+1}$, and conversely.

```
In[31]:= SubstTest[member, nat[t], natmul[nat[t], nat[w]],
  t -> natexp[nat[w], nat[x]]] // MapNotNot // Reverse
```

```
Out[31]= member[natexp[nat[w], nat[x]], natmul[nat[w], natexp[nat[w], nat[x]]]] ==
  member[set[0], nat[w]]
```

```
In[32]:= member[natexp[nat[w_], nat[x_]], natmul[nat[w_], natexp[nat[w_], nat[x_]]]] :=
  member[set[0], nat[w]]
```

A form of the transitive law is needed: if $b^m < b^{m+1}$ and $b^{m+1} \leq b^n$, then $b^m < b^n$.

```
In[33]:= SubstTest[implies,
  and[member[nat[t], nat[u]], subclass[nat[u], nat[v]], member[nat[t], nat[v]],
  {t -> natexp[nat[w], nat[x]], u -> natexp[nat[w], succ[nat[x]]],
  v -> natexp[nat[w], nat[y]]}] // Reverse
```

```
Out[33]= or[member[natexp[nat[w], nat[x]], natexp[nat[w], nat[y]]],
  member[natexp[nat[w], nat[y]], natmul[nat[w], natexp[nat[w], nat[x]]]],
  not[member[set[0], nat[w]]]] == True
```

```
In[34]:= (% /. {w -> w_, x -> x_, y -> y_}) /. Equal -> SetDelayed
```

Main theorem.

```
In[35]:= Map[not, SubstTest[and, implies[p1, p2], implies[and[p1, p4], p5],
  implies[and[p2, p3], p4], not[implies[and[p1, p3], p5]],
  {p1 -> member[set[0], nat[w]], p2 -> not[equal[0, nat[w]]], p3 -> member[nat[x], nat[y]],
  p4 -> not[member[natexp[nat[w], nat[y]], natmul[nat[w], natexp[nat[w], nat[x]]]]],
  p5 -> member[natexp[nat[w], nat[x]], natexp[nat[w], nat[y]]]}] // Reverse
```

```
Out[35]= or[member[natexp[nat[w], nat[x]], natexp[nat[w], nat[y]]],
  not[member[nat[x], nat[y]], not[member[set[0], nat[w]]]] == True
```

```
In[36]:= or[member[natexp[nat[w_], nat[x_]], natexp[nat[w_], nat[y_]]],
  not[member[nat[x_], nat[y_]], not[member[set[0], nat[w_]]]] := True
```

The hypothesis `member[set[0],nat[w]]` can be replaced with a double `succ` wrapper. A natural number is greater than `1` if and only if it is the successor of the successor of some number.

```
In[37]:= SubstTest[implies, and[member[set[0], nat[t]], member[nat[x], nat[y]]],
  member[natexp[nat[t], nat[x]], natexp[nat[t], nat[y]]],
  t -> succ[succ[nat[w]]] // Reverse
```

```
Out[37]= or[member[natexp[succ[succ[nat[w]]], nat[x]], natexp[succ[succ[nat[w]]], nat[y]]],
  not[member[nat[x], nat[y]]]] == True
```

```
In[38]:= (% /. {w -> w_, x -> x_, y -> y_}) /. Equal -> SetDelayed
```

The converse also holds:

```
In[39]:= SubstTest[implies, member[natexp[succ[nat[t]], nat[x]], natexp[succ[nat[t]], nat[y]]],
  member[nat[x], nat[y]], t → succ[nat[w]]] // Reverse
```

```
Out[39]= or[member[nat[x], nat[y]], not[member[
  natexp[succ[succ[nat[w]]], nat[x]], natexp[succ[succ[nat[w]]], nat[y]]]]] = True
```

```
In[40]:= (% /. {w → w_, x → x_, y → y_}) /. Equal → SetDelayed
```

Combining the two directions, one finds:

```
In[41]:= equiv[member[natexp[succ[succ[nat[w]]], nat[x]], natexp[succ[succ[nat[w]]], nat[y]]],
  member[nat[x], nat[y]]]
```

```
Out[41]= True
```

```
In[42]:= member[natexp[succ[succ[nat[w_]]], nat[x_]], natexp[succ[succ[nat[w_]]], nat[y_]]] :=
  member[nat[x], nat[y]]
```

Comment. The special case $w = 0$ was derived previously using a cardinality argument.