

# PRELUDE

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```
In[1]:= SetDirectory["1:"]; << goedel.12may15a

:Package Title: goedel.12may15a                2012 May 15 at 1:05 p.m.

Loading takes about seventeen minutes, half that time due to builtin pauses.

It is now: 2012 May 16 at 9:11

Loading Simplification Rules

TOOLS.M is now incorporated in the GOEDEL program as of 2010 September 3

weightlimit = 40

Loading completed.

It is now: 2012 May 16 at 9:28
```

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## summary

In the monograph *On numbers and games* John H. Conway in 1976 showed that the class of ordinals can be embedded in a (proper class) field. Philip Ehrlich simplified Conway's constructions by introducing a relation here called **PRELUDE**.

```
In[2]:= "Philip Ehrlich, The Absolute Arithmetic  
Continuum and the Unification of all Numbers Great and Small  
The Bulletin of Symbolic Logic, volume 18, pages 145 (2012)";
```

The intention is to use this relation to apply the general theory of well-founded recursion to define important functions in the theory of Conway games. In this notebook this relation is formally defined and it is shown that it is well-founded and that its inverse is thin.

---

## review

A brief review of the development of the theory of Conway games within the **GOEDEL** program will suffice. The class **GAMES** of Conway games can be characterized as the smallest class satisfying  $\mathbf{P[x]} \times \mathbf{P[x]} \subset \mathbf{x}$ . The existence of this class was established in a series of three notebooks using a transfinite well-founded recursive construction involving the function **gamerec**, defined as follows.

```
In[3]:= composite[rec[composite[CART, DUP, POWER, BIGCUP, IMAGE, SECOND], SECOND],
               composite[inverse[E], id[OMEGA]]], id[OMEGA]]
```

```
Out[3]= gamerec
```

The class **GAMES** is the sum class of the range of the function **gamerec**.

```
In[4]:= U[range[gamerec]]
```

```
Out[4]= GAMES
```

The class **GAMES** was shown to satisfy the following three rewrite rules:

```
In[5]:= fix[GAMES]
```

```
Out[5]= P[GAMES]
```

```
In[6]:= cart[P[GAMES], P[GAMES]]
```

```
Out[6]= GAMES
```

```
In[7]:= implies[subclass[cart[P[x], P[x]], x], subclass[GAMES, x]]
```

```
Out[7]= True
```

## the birthday function

A game **x** is said to be **alive** at time **t** if **t** is an ordinal, and  $x \in \text{APPLY}[\text{gamerec}, t]$ . The **birthday** of a game **x** is the least ordinal **t** for which **x** is alive at time **t**. The birthday function **BDAY** is defined by the rewrite rule:

```
In[8]:= VERTSECT[complement[composite[complement[inverse[E]], inverse[gamerec], E]]]
```

```
Out[8]= BDAY
```

A game **x** is alive at any (ordinal) time **ord[t]** except before its birth.

```
In[9]:= member[ord[t], APPLY[BDAY, x]]
```

```
Out[9]= not[member[x, APPLY[gamerec, ord[t]]]]
```

Theorem. The relation  $\text{inverse}[\text{BDAY}] \circ E \circ \text{BDAY}$  is well-founded.

```
In[10]:= SubstTest[WELLFOUNDED, composite[inverse[funpart[u]], wf[v], funpart[u]],
                 {u → BDAY, v → composite[id[OMEGA], E]}] // Reverse
```

```
Out[10]= WELLFOUNDED[composite[inverse[BDAY], E, BDAY]] == True
```

```
In[11]:= WELLFOUNDED[composite[inverse[BDAY], E, BDAY]] := True
```

---

## definition of PRELUDE

Any (Conway) game is an ordered pair of sets of games, called the **options** for the players "left" and "right". It will be said that a game  $x$  is a **prelude** to a game  $y$  if  $x$  is one of the left or right options of the game  $y$ . Formally the relation **PRELUDE** is defined as follows.

```
In[12]:= union[composite[id[GAMES], inverse[FIRST], E],
           composite[id[GAMES], inverse[SECOND], E]] := PRELUDE
```

Theorem. Domain of the **PRELUDE** relation.

```
In[13]:= SubstTest[domain, union[u, v], {u -> composite[id[GAMES], inverse[FIRST], E],
           v -> composite[id[GAMES], inverse[SECOND], E]}] // Reverse
```

```
Out[13]= domain[PRELUDE] == GAMES
```

```
In[14]:= domain[PRELUDE] := GAMES
```

Theorem. An upper bound for the range of the **PRELUDE** relation.

```
In[15]:= Map[subclass[#, GAMES] &,
           SubstTest[range, union[u, v], {u -> composite[id[GAMES], inverse[FIRST], E],
           v -> composite[id[GAMES], inverse[SECOND], E]}]] // Reverse
```

```
Out[15]= subclass[range[PRELUDE], GAMES] == True
```

```
In[16]:= subclass[range[PRELUDE], GAMES] := True
```

---

## simplification rules for PRELUDE

Theorem. Simplification rule.

```
In[17]:= SubstTest[composite, Id, union[u, v], {u -> composite[id[GAMES], inverse[FIRST], E],
           v -> composite[id[GAMES], inverse[SECOND], E]}] // Reverse
```

```
Out[17]= composite[Id, PRELUDE] == PRELUDE
```

```
In[18]:= composite[Id, PRELUDE] := PRELUDE
```

Theorem. Normalization rule for the inverse of the **PRELUDE** relation.

```
In[19]:= SubstTest[inverse, union[u, v], {u -> composite[id[GAMES], inverse[FIRST], E],
           v -> composite[id[GAMES], inverse[SECOND], E]}]
```

```
Out[19]= union[composite[inverse[E], FIRST, id[GAMES]],
              composite[inverse[E], SECOND, id[GAMES]]] == inverse[PRELUDE]
```

```
In[20]:= union[composite[inverse[E], FIRST, id[GAMES]],
             composite[inverse[E], SECOND, id[GAMES]]] := inverse[PRELUDE]
```

Theorem. Simplification rule.

```
In[21]:= SubstTest[composite, id[GAMES],
                 union[u, v], {u -> composite[id[GAMES], inverse[FIRST], E],
                               v -> composite[id[GAMES], inverse[SECOND], E]}] // Reverse
```

```
Out[21]= composite[id[GAMES], PRELUDE] == PRELUDE
```

```
In[22]:= composite[id[GAMES], PRELUDE] := PRELUDE
```

Theorem. Simplification rule.

```
In[23]:= ImageComp[id[GAMES], PRELUDE, x] // Reverse
```

```
Out[23]= intersection[GAMES, image[PRELUDE, x]] == image[PRELUDE, x]
```

```
In[24]:= intersection[GAMES, image[PRELUDE, x_]] := image[PRELUDE, x]
```

Theorem. Simplification rule.

```
In[25]:= IminComp[id[GAMES], PRELUDE, x] // Reverse
```

```
Out[25]= image[inverse[PRELUDE], intersection[GAMES, x]] == image[inverse[PRELUDE], x]
```

```
In[26]:= image[inverse[PRELUDE], intersection[GAMES, x_]] := image[inverse[PRELUDE], x]
```

Theorem. Simplification rule.

```
In[27]:= ImageComp[PRELUDE, id[GAMES], x] // Reverse
```

```
Out[27]= image[PRELUDE, intersection[GAMES, x]] == image[PRELUDE, x]
```

```
In[28]:= image[PRELUDE, intersection[GAMES, x_]] := image[PRELUDE, x]
```

Theorem. Simplification rule.

```
In[29]:= IminComp[PRELUDE, id[GAMES], x] // Reverse
```

```
Out[29]= intersection[GAMES, image[inverse[PRELUDE], x]] == image[inverse[PRELUDE], x]
```

```
In[30]:= intersection[GAMES, image[inverse[PRELUDE], x_]] := image[inverse[PRELUDE], x]
```

Lemma. Simplification rule.

```
In[31]:= equiv[and[member[x, GAMES], member[first[x], V]], member[x, GAMES]]
```

```
Out[31]= True
```

```
In[32]:= and[member[x_, GAMES], member[first[x_], V]] := member[x, GAMES]
```

Theorem. Membership rule.

```

In[33]:= SubstTest[member, pair[x, y], union[u, v], {u -> composite[id[GAMES], inverse[FIRST], E],
  v -> composite[id[GAMES], inverse[SECOND], E]}] // Reverse

Out[33]= member[pair[x, y], PRELUDE] == or[and[member[x, first[y]], member[y, GAMES]],
  and[member[x, second[y]], member[y, GAMES]]]

In[34]:= member[pair[x_, y_], PRELUDE] := or[and[member[x, first[y]], member[y, GAMES]],
  and[member[x, second[y]], member[y, GAMES]]]

```

---

## birthdays of preludes

It is shown in this section that the preludes (left and right options) of a game were born before the game itself.

Theorem. If a game  $x$  is alive at time  $t$ , then  $\text{first}[x] \subset U[\text{image}[\text{gamerec}, t]]$ .

```

In[35]:= Map[implies[#, subclass[first[x], U[image[gamerec, ord[t]]]]] &,
  SubstTest[member, x, cartsq[w], w -> P[U[image[gamerec, ord[t]]]]] // Reverse

Out[35]= or[not[member[x, APPLY[gamerec, ord[t]]]],
  subclass[first[x], U[image[gamerec, ord[t]]]] == True

In[36]:= or[not[member[x_, APPLY[gamerec, ord[t_]]]],
  subclass[first[x_], U[image[gamerec, ord[t_]]]] := True

```

Dual. If a game  $x$  is alive at time  $t$ , then  $\text{second}[x] \subset U[\text{image}[\text{gamerec}, t]]$ .

```

In[37]:= Map[implies[#, subclass[second[x], U[image[gamerec, ord[t]]]]] &,
  SubstTest[member, x, cartsq[w], w -> P[U[image[gamerec, ord[t]]]]] // Reverse

Out[37]= or[not[member[x, APPLY[gamerec, ord[t]]]],
  subclass[second[x], U[image[gamerec, ord[t]]]] == True

In[38]:= or[not[member[x_, APPLY[gamerec, ord[t_]]]],
  subclass[second[x_], U[image[gamerec, ord[t_]]]] := True

```

Corollary. (The **ord** wrapper is removed, and the above result is then specialized to the birthday of the game. This has a redundant literal.)

```

In[39]:= (SubstTest[implies, equal[t, ord[w]], or[not[member[x, APPLY[gamerec, t]]],
  subclass[first[x], U[image[gamerec, t]]], w -> t] /. t -> APPLY[BDAY, x]) // Reverse

Out[39]= or[not[member[x, GAMES]], not[member[x, APPLY[gamerec, APPLY[BDAY, x]]]],
  subclass[first[x], U[image[gamerec, APPLY[BDAY, x]]]] == True

In[40]:= (% /. x -> x_) /. Equal -> SetDelayed

```

The following result is better.

Theorem.

```
In[41]:= Map[not, SubstTest[and, implies[p1, p2], implies[and[p1, p2], p3], not[implies[p1, p3]],
  {p1 -> member[x, GAMES], p2 -> member[x, APPLY[gamerec, APPLY[BDAY, x]]],
  p3 -> subclass[first[x], U[image[gamerec, APPLY[BDAY, x]]]}] // Reverse
```

```
Out[41]= or[not[member[x, GAMES]], subclass[first[x], U[image[gamerec, APPLY[BDAY, x]]]] = True
```

```
In[42]:= or[not[member[x_, GAMES]],
  subclass[first[x_], U[image[gamerec, APPLY[BDAY, x_]]]] := True
```

Lemma for the dual.

```
In[43]:= (SubstTest[implies, equal[t, ord[w]], or[not[member[x, APPLY[gamerec, t]]],
  subclass[second[x], U[image[gamerec, t]]], w -> t] /. t -> APPLY[BDAY, x]) // Reverse
```

```
Out[43]= or[not[member[x, GAMES]], not[member[x, APPLY[gamerec, APPLY[BDAY, x]]],
  subclass[second[x], U[image[gamerec, APPLY[BDAY, x]]]] = True
```

```
In[44]:= (% /. x -> x_) /. Equal -> SetDelayed
```

Dual Theorem.

```
In[45]:= Map[not, SubstTest[and, implies[p1, p2], implies[and[p1, p2], p3], not[implies[p1, p3]],
  {p1 -> member[x, GAMES], p2 -> member[x, APPLY[gamerec, APPLY[BDAY, x]]],
  p3 -> subclass[second[x], U[image[gamerec, APPLY[BDAY, x]]]}] // Reverse
```

```
Out[45]= or[not[member[x, GAMES]], subclass[second[x], U[image[gamerec, APPLY[BDAY, x]]]] = True
```

```
In[46]:= or[not[member[x_, GAMES]],
  subclass[second[x_], U[image[gamerec, APPLY[BDAY, x_]]]] := True
```

Lemma.

```
In[47]:= SubstTest[implies, and[subclass[u, v], subclass[v, w]], subclass[u, w],
  {u -> range[BDAY], v -> OMEGA, w -> union[image[S, set[ord[t]]], ord[t]]} // Reverse
```

```
Out[47]= subclass[range[BDAY], union[image[S, set[ord[t]]], ord[t]] = True
```

```
In[48]:= (% /. t -> t_) /. Equal -> SetDelayed
```

Theorem.

```
In[49]:= Map[subclass[image[#, ord[t]], ord[t]] &,
  Assoc[BDAY, inverse[BDAY], composite[complement[E], id[OMEGA]]]
```

```
Out[49]= subclass[image[BDAY, U[image[gamerec, ord[t]]], ord[t]] = True
```

```
In[50]:= subclass[image[BDAY, U[image[gamerec, ord[t_]]], ord[t_]] := True
```

Lemma. (The **ord** wrapper is removed, and the result is then specialized to the birthday of the game.)

```
In[51]:= (SubstTest[implies, equal[t, ord[w]], subclass[image[BDAY, U[image[gamerec, t]]], t],
          w → t] // Reverse) /. t → APPLY[BDAY, x]
```

```
Out[51]= or[not[member[x, GAMES]],
           subclass[image[BDAY, U[image[gamerec, APPLY[BDAY, x]]]], APPLY[BDAY, x]]] == True
```

```
In[52]:= (% /. x → x_) /. Equal → SetDelayed
```

Lemma.

```
In[53]:= Map[not, SubstTest[and, implies[p1, p2], implies[p2, p3], implies[p1, p4],
                        implies[and[p3, p4], p5], not[implies[p1, p5]], {p1 → member[x, GAMES],
                        p2 → subclass[first[x], U[image[gamerec, APPLY[BDAY, x]]]}, p3 →
                        subclass[image[BDAY, first[x]], image[BDAY, U[image[gamerec, APPLY[BDAY, x]]]}],
                        p4 → subclass[image[BDAY, U[image[gamerec, APPLY[BDAY, x]]]], APPLY[BDAY, x]],
                        p5 → subclass[image[BDAY, first[x]], APPLY[BDAY, x]]}] // Reverse
```

```
Out[53]= or[not[member[x, GAMES]], subclass[image[BDAY, first[x]], APPLY[BDAY, x]]] == True
```

```
In[54]:= (% /. x → x_) /. Equal → SetDelayed
```

Dual lemma.

```
In[55]:= Map[not, SubstTest[and, implies[p1, p2], implies[p2, p3], implies[p1, p4],
                        implies[and[p3, p4], p5], not[implies[p1, p5]], {p1 → member[x, GAMES],
                        p2 → subclass[second[x], U[image[gamerec, APPLY[BDAY, x]]]}, p3 →
                        subclass[image[BDAY, second[x]], image[BDAY, U[image[gamerec, APPLY[BDAY, x]]]}],
                        p4 → subclass[image[BDAY, U[image[gamerec, APPLY[BDAY, x]]]], APPLY[BDAY, x]],
                        p5 → subclass[image[BDAY, second[x]], APPLY[BDAY, x]]}] // Reverse
```

```
Out[55]= or[not[member[x, GAMES]], subclass[image[BDAY, second[x]], APPLY[BDAY, x]]] == True
```

```
In[56]:= (% /. x → x_) /. Equal → SetDelayed
```

The preceding two lemmas both contain a redundant literal. The following result allows one to eliminate this literal.

Lemma.

```
In[57]:= SubstTest[implies, equal[v, V], subclass[u, v], v → APPLY[BDAY, x]] // Reverse
```

```
Out[57]= or[member[x, GAMES], subclass[u, APPLY[BDAY, x]]] == True
```

```
In[58]:= (% /. {x → x_, u → u_}) /. Equal → SetDelayed
```

Theorem. The left options of a game were born before the game itself.

```
In[59]:= SubstTest[and, implies[p, q], or[p, q],
                {p → member[x, GAMES], q → subclass[image[BDAY, first[x]], APPLY[BDAY, x]]}]
```

```
Out[59]= subclass[image[BDAY, first[x]], APPLY[BDAY, x]] == True
```

```
In[60]:= subclass[image[BDAY, first[x_]], APPLY[BDAY, x_]] := True
```

Dual Theorem. The right options of a game were born before the game itself.

```
In[61]:= SubstTest[and, implies[p, q], or[p, q],
  {p -> member[x, GAMES], q -> subclass[image[BDAY, second[x]], APPLY[BDAY, x]]}]
```

```
Out[61]= subclass[image[BDAY, second[x]], APPLY[BDAY, x]] == True
```

```
In[62]:= subclass[image[BDAY, second[x_]], APPLY[BDAY, x_]] := True
```

The next step is to eliminate the variable  $x$ .

Theorem.

```
In[63]:= Map[equal[V, domain[#]] &,
  SubstTest[reify, x, case[subclass[image[t, first[x]], APPLY[t, x]], t -> BDAY]]
```

```
Out[63]= subclass[GAMES, fix[composite[inverse[BDAY], S, IMAGE[BDAY], FIRST]]] == True
```

```
In[64]:= subclass[GAMES, fix[composite[inverse[BDAY], S, IMAGE[BDAY], FIRST]]] := True
```

Dual theorem.

```
In[65]:= Map[equal[V, domain[#]] &,
  SubstTest[reify, x, case[subclass[image[t, second[x]], APPLY[t, x]], t -> BDAY]]
```

```
Out[65]= subclass[GAMES, fix[composite[inverse[BDAY], S, IMAGE[BDAY], SECOND]]] == True
```

```
In[66]:= subclass[GAMES, fix[composite[inverse[BDAY], S, IMAGE[BDAY], SECOND]]] := True
```

Corollary.

```
In[67]:= subclass[composite[id[GAMES], inverse[FIRST], E],
  composite[inverse[BDAY], E, BDAY]] // AssertTest // InvertFix
```

```
Out[67]= subclass[composite[id[GAMES], inverse[FIRST], E],
  composite[inverse[BDAY], E, BDAY]] == True
```

```
In[68]:= subclass[composite[id[GAMES], inverse[FIRST], E],
  composite[inverse[BDAY], E, BDAY]] := True
```

Dual corollary.

```
In[69]:= subclass[composite[id[GAMES], inverse[SECOND], E],
  composite[inverse[BDAY], E, BDAY]] // AssertTest // InvertFix
```

```
Out[69]= subclass[composite[id[GAMES], inverse[SECOND], E],
  composite[inverse[BDAY], E, BDAY]] == True
```

```
In[70]:= subclass[composite[id[GAMES], inverse[SECOND], E],
  composite[inverse[BDAY], E, BDAY]] := True
```

The preceding two results can be combined as follows.



Theorem. The relation **PRELUDE** is a subclass of the well-founded relation  $\text{inverse}[\text{BDAY}] \circ \text{E} \circ \text{BDAY}$ .

```
In[71]:= SubstTest[subclass, union[u, v], composite[inverse[BDAY], E, BDAY],
  {u -> composite[id[GAMES], inverse[FIRST], E],
   v -> composite[id[GAMES], inverse[SECOND], E]}] // Reverse
```

```
Out[71]= subclass[PRELUDE, composite[inverse[BDAY], E, BDAY]] == True
```

```
In[72]:= subclass[PRELUDE, composite[inverse[BDAY], E, BDAY]] := True
```

Theorem. The relation **PRELUDE** is well-founded.

```
In[73]:= SubstTest[implies, and[subclass[u, v], WELLFOUNDED[v]], WELLFOUNDED[u],
  {u -> PRELUDE, v -> composite[inverse[BDAY], E, BDAY]}] // Reverse
```

```
Out[73]= WELLFOUNDED[PRELUDE] == True
```

```
In[74]:= WELLFOUNDED[PRELUDE] := True
```

---

## VERTSECT rule

Theorem.

```
In[75]:= SubstTest[VERTSECT, union[u, v], {u -> inverse[composite[id[GAMES], inverse[FIRST], E]],
  v -> inverse[composite[id[GAMES], inverse[SECOND], E]]}] // Reverse
```

```
Out[75]= VERTSECT[inverse[PRELUDE]] ==
  union[cart[complement[GAMES], set[0]], composite[CUP, id[GAMES]]]
```

```
In[76]:= VERTSECT[inverse[PRELUDE]] :=
  union[cart[complement[GAMES], set[0]], composite[CUP, id[GAMES]]]
```

Corollary. The inverse of **PRELUDE** is thin. No new rewrite rule is needed for this.

```
In[77]:= thin[inverse[PRELUDE]]
```

```
Out[77]= True
```