

## two-sided restrictions to final segments

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```
In[1]:= SetDirectory["1:"]; << goedel.10jul13a; << tools.m

:Package Title: goedel.10jul13a          2010 July 13 at 12:20 p.m.

It is now: 2010 Jul 15 at 14:45

Loading Simplification Rules

TOOLS.M                                Revised 2010 February 26

weightlimit = 40
```

---

### summary

In the recently posted notebook **trv-fnsg.nb** a variable-free statement about transitivity for final segments was derived. In the present notebook another variable-free statement concerning the transitivity of final segments is derived in which the predicate **TRANSITIVE** occurs explicitly.

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### acknowledgment

Theorem 8.10 on page 32 in the following reference inspired this work.

```
In[2]:= "Egbert Harzheim, Ordered Sets, Advances in Mathematics, volume 7, Springer
        Science+Business Media, Inc., 2005. ISBN 0387-24219-8. QA171.48 .H37";
```

---

### temporary abbreviations

Two-sided restriction can be viewed as an action of sets on relations. If  $s$  is a set and  $r$  is a relation, then this action of  $s$  on  $r$  could be denoted by  $s \cdot r = \text{restrict}[r, s, s]$ . Actually one can allow  $r$  to be any set. There is no need to require  $r$  to be a relation. The following temporary abbreviation for the binary function for this action will be used.

```
In[3]:= TWOSIDED := composite[CAP, cross[composite[CART, DUP], Id]]
```

The **APPLY** rule for this action simplifies when one uses **setpart** wrappers:

```
In[4]:= APPLY[TWOSIDED, PAIR[setpart[s], setpart[r]]]
```

```
Out[4]= composite[id[setpart[s]], setpart[r], id[setpart[s]]]
```

Harzheim defines a set  $s$  to be a **final segment** of a relation  $r$  if  $\text{image}[r, s] = s$ . The following temporary abbreviation for the relation of being a final segment will be used here.

```
In[5]:= FINSEG := fix[composite[inverse[SECOND], IMG]]
```

The following membership rule for this relation helps explain its significance.

```
In[6]:= member[pair[setpart[r], setpart[s]], FINSEG]
```

```
Out[6]= equal[image[setpart[r], setpart[s]], setpart[s]]
```

This notebook concerns only the restriction of the two-sided restriction action to final segments:

```
In[7]:= FNRS := composite[TWOSIDED, id[inverse[FINSEG]]]
```

It will be shown below that the relation  $\text{FNRS} \circ \text{inverse}[\text{SECOND}]$  is transitive. An ordered pair of relations  $\text{pair}[\mathbf{x}, \mathbf{y}]$  belongs to this transitive relation if  $\mathbf{y}$  is the two-sided restriction of  $\mathbf{x}$  to some final segment of  $\mathbf{x}$ .

## a quasi-associative law

The **TWOSIDED** action satisfies the following quasi-associative law:  $\mathbf{u} \cdot (\mathbf{v} \cdot \mathbf{x}) = (\mathbf{u} \cap \mathbf{v}) \cdot \mathbf{x}$ . This quasi-associative law simplifies for the restricted action **FNRS** in that  $\mathbf{u} \cap \mathbf{v}$  can be replaced with  $\mathbf{u}$ . Explicitly, if one writes  $\mathbf{r} \cdot \mathbf{s}$  for this restriction of the action  $\mathbf{r} \cdot \mathbf{s}$ , the quasi-associative law effectively becomes the inclusion  $\mathbf{u} \cdot (\mathbf{v} \cdot \mathbf{x}) \subset \mathbf{u} \cdot \mathbf{x}$ . (Note that one must replace equality by inclusion here because the restricted action is no longer total. The following lemma formally expresses this simplification of the quasi-associative law for the action **FNRS** in the form of an implication involving two additional variables:  $\mathbf{y} = \mathbf{v} \cdot \mathbf{x} \ \& \ \mathbf{z} = \mathbf{u} \cdot \mathbf{y} \implies \mathbf{z} = \mathbf{u} \cdot \mathbf{x}$ .)

Lemma. A quasi-associative law for the action **FNRS** with five variables wrapped with **setpart**.

```
In[8]:= implies[and[member[pair[pair[setpart[u], setpart[x]], setpart[y]], FNRS],
  member[pair[pair[setpart[v], setpart[y]], setpart[z]], FNRS]],
  member[pair[pair[setpart[v], setpart[x]], setpart[z]], FNRS]] // NotNotTest
```

```
Out[8]= or[and[equal[composite[id[setpart[v]], setpart[x], id[setpart[v]]], setpart[z]],
  equal[image[setpart[x], setpart[v]], setpart[v]]],
  not[equal[composite[id[setpart[u]], setpart[x], id[setpart[u]]], setpart[y]]],
  not[equal[composite[id[setpart[v]], setpart[y], id[setpart[v]]], setpart[z]]],
  not[equal[image[setpart[x], setpart[u]], setpart[u]]],
  not[equal[image[setpart[y], setpart[v]], setpart[v]]]] = True
```

```
In[9]:= (% /. {u -> u_, v -> v_, x -> x_, y -> y_, z -> z_}) /. Equal -> SetDelayed
```

The main problem encountered in eliminating the variables is that the complexity of the formula for **FNRS** causes various rewrite rules to produce complicated expressions. To shield **FNRS** from unwanted rewriting in the process of eliminating variables, yet another variable  $\mathbf{t}$  and an equality literal will be introduced. Accordingly, the above simplified statement of the quasi-associative law will first be rewritten as follows:

```
In[10]:= implies[and[equal[t, FNRS], member[pair[pair[setpart[v], setpart[x]], setpart[y]], t],
  member[pair[pair[setpart[u], setpart[y]], setpart[z]], t]],
  member[pair[pair[setpart[u], setpart[x]], setpart[z]], t]]
```

```
Out[10]= True
```

The goal is to eliminate all six variables to derive a variable-free statement of transitivity.

## eliminating x, y and z

Lemma. (Elimination of three variables.)

```
In[11]:= Map[empty[composite[complement[#], id[cart[V, V]]]] &,
  SubstTest[class, pair[pair[x, z], y],
    implies[and[equal[t, s], member[pair[pair[setpart[v], setpart[x]], setpart[y]], t],
      member[pair[pair[setpart[u], setpart[y]], setpart[z]], t]],
      member[pair[pair[setpart[u], setpart[x]], setpart[z]], t]], s → FNRS]]
```

```
Out[11]= or[not[equal[0, fix[composite[fix[composite[inverse[SECOND], IMG]], fix[
  composite[inverse[t], complement[CAP], cross[composite[CART, DUP], Id]]]]]],
  not[subclass[t, cart[inverse[fix[composite[inverse[SECOND], IMG]], V]]],
  not[subclass[inverse[fix[composite[inverse[SECOND], IMG]]],
    fix[composite[inverse[t], CAP, cross[composite[CART, DUP], Id]]]],
  subclass[composite[t, LEFT[setpart[u]], t, LEFT[setpart[v]],
    composite[t, LEFT[setpart[u]]]]] == True
```

```
In[12]:= (% /. {t → t_, u → u_, v → v_}) /. Equal → SetDelayed
```

This rather messy statement can be cleaned up as follows:

Theorem. The composite of the functions  $\text{FNRS} \circ \text{LEFT}[u]$  and  $\text{FNRS} \circ \text{LEFT}[v]$  is contained in the left hand factor:  $\text{FNRS} \circ \text{LEFT}[u]$ .

```
In[13]:= Map[not, SubstTest[and, implies[p1, p2], implies[p1, p3], implies[p1, p4],
  implies[and[p2, p3, p4], p5], not[implies[p1, p5]], {p1 → equal[t, FNRS],
  p2 → equal[0, fix[composite[fix[composite[inverse[SECOND], IMG]], fix[
    composite[inverse[t], complement[CAP], cross[composite[CART, DUP], Id]]]]]],
  p3 → subclass[t, cart[inverse[fix[composite[inverse[SECOND], IMG]], V]],
  p4 → subclass[inverse[fix[composite[inverse[SECOND], IMG]]],
    fix[composite[inverse[t], CAP, cross[composite[CART, DUP], Id]]]],
  p5 → subclass[composite[t, LEFT[setpart[u]], t, LEFT[setpart[v]],
    composite[t, LEFT[setpart[u]]]]]] // Reverse
```

```
Out[13]= or[not[equal[t, composite[CAP, cross[composite[CART, DUP], Id],
  id[inverse[fix[composite[inverse[SECOND], IMG]]]]]],
  subclass[composite[t, LEFT[setpart[u]], t, LEFT[setpart[v]],
    composite[t, LEFT[setpart[u]]]]] == True
```

```
In[14]:= (% /. {t → t_, u → u_, v → v_}) /. Equal → SetDelayed
```

The following observation implies a corollary.

```
In[15]:= subclass[LEFT[u], inverse[SECOND]]
```

```
Out[15]= True
```

Corollary. The composite of the functions  $\mathbf{FNRS} \circ \mathbf{LEFT}[u]$  and  $\mathbf{FNRS} \circ \mathbf{LEFT}[v]$  is contained in the relation  $\mathbf{FNRS} \circ \mathbf{inverse}[\mathbf{SECOND}]$ .

```
In[16]:= Map[not, SubstTest[and, implies[p1, p2],
  implies[p2, p3], not[implies[p1, p3]], {p1 → equal[t, FNRS],
  p2 → subclass[composite[t, LEFT[setpart[u]], t, LEFT[setpart[v]]],
  composite[t, LEFT[setpart[u]]], p3 → subclass[composite[t, LEFT[setpart[u]],
  t, LEFT[setpart[v]]], composite[t, inverse[SECOND]]}]] // Reverse
```

```
Out[16]= or[not[equal[t, composite[CAP, cross[composite[CART, DUP], Id],
  id[inverse[fix[composite[inverse[SECOND], IMG]]]]]],
  subclass[composite[t, LEFT[setpart[u]], t, LEFT[setpart[v]]],
  composite[t, inverse[SECOND]]] == True
```

```
In[17]:= (% /. {t → t_, u → u_, v → v_}) /. Equal → SetDelayed
```

After some experimentation it was discovered that it is much easier to eliminate the variables  $u$  and  $v$  if one works with right multiplications for the flipped action  $\mathbf{flip}[\mathbf{FNRS}]$  instead of left multiplications of  $\mathbf{FNRS}$ . Some additional steps are needed to do this.

Corollary.

```
In[18]:= SubstTest[implies, equal[s, FNRS],
  subclass[composite[s, LEFT[setpart[u]], s, LEFT[setpart[v]]],
  composite[s, inverse[SECOND]]], s → flip[t] // Reverse
```

```
Out[18]= or[not[equal[composite[t, SWAP], composite[CAP, cross[composite[CART, DUP], Id],
  id[inverse[fix[composite[inverse[SECOND], IMG]]]]]],
  subclass[composite[t, RIGHT[setpart[u]], t, RIGHT[setpart[v]]],
  composite[t, inverse[FIRST]]] == True
```

```
In[19]:= (% /. {t → t_, u → u_, v → v_}) /. Equal → SetDelayed
```

Corollary.

```
In[20]:= (Map[not, SubstTest[and, implies[and[p1, p2], p3],
  implies[and[p2, p3], p4], not[implies[and[p1, p2], p4]],
  {p1 → equal[s, flip[FNRS]], p2 → equal[s, t], p3 → equal[flip[t], FNRS],
  p4 → subclass[composite[t, RIGHT[setpart[u]], t, RIGHT[setpart[v]]],
  composite[t, inverse[FIRST]]}]] // Reverse) /. s → flip[FNRS]
```

```
Out[20]= or[not[equal[t, composite[CAP,
  cross[Id, composite[CART, DUP]], id[fix[composite[inverse[SECOND], IMG]]]]],
  subclass[composite[t, RIGHT[setpart[u]], t, RIGHT[setpart[v]]],
  composite[t, inverse[FIRST]]] == True
```

```
In[21]:= (% /. {t → t_, u → u_, v → v_}) /. Equal → SetDelayed
```

All the remaining variables can now be eliminated all at once to obtain a transitive law.

Theorem. Two-sided restriction to final segments is transitive.

```
In[22]:= Map[empty[composite[Id, complement[#]]] &, SubstTest[class, pair[u, v],
  implies[equal[t, s], subclass[composite[t, RIGHT[setpart[u]], t, RIGHT[setpart[v]]],
    composite[t, inverse[FIRST]]], s → flip[FNRS]]] /. t → flip[FNRS]
```

```
Out[22]= TRANSITIVE[
  composite[CAP, id[composite[CART, DUP, fix[composite[inverse[SECOND], IMG]]]],
    inverse[FIRST]]] = True
```

```
In[23]:= TRANSITIVE[
  composite[CAP, id[composite[CART, DUP, fix[composite[inverse[SECOND], IMG]]]],
    inverse[FIRST]]] := True
```

Restatement.

```
In[24]:= TRANSITIVE[composite[FNRS, inverse[SECOND]]]
```

```
Out[24]= True
```

## final comments

Comment. The relation  $\text{FNRS} \circ \text{inverse}[\text{SECOND}]$  is also antisymmetric.

```
In[25]:= ANTISYMMETRIC[composite[FNRS, inverse[SECOND]]]
```

```
Out[25]= True
```

The union of the domain and range of  $\text{FNRS} \circ \text{inverse}[\text{SECOND}]$  is the universal class  $V$ .

```
In[33]:= udora[composite[FNRS, inverse[SECOND]]]
```

```
Out[33]= V
```

The following statement also follows automatically now.

```
In[36]:= PARTIALORDER[union[Id, composite[FNRS, inverse[SECOND]]]]
```

```
Out[36]= True
```