

a strong form of reflexivity

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```
In[1]:= SetDirectory["1:"]; << goedel.08feb10a; << tools.m

:Package Title: goedel.08feb10a          2008 February 10 at 2:15 p.m.

It is now: 2008 Feb 12 at 16:40

Loading Simplification Rules

TOOLS.M                                Revised 2008 January 22

weightlimit = 40
```

summary

This notebook is concerned with a careful study of the concept of a reflexive relation. Patrick Suppes introduced the following definition for a strong form of reflexivity of a relation x on a set y . (Actually, both x and y can be arbitrary classes here.)

```
In[2]:= assert[forall[t, implies[member[t, y], member[pair[t, t], x]]]]
Out[2]= subclass[y, fix[x]]
```

Reference: (see Definition 10 on page 69).

```
In[3]:= "Patrick Suppes, Axiomatic Set Theory, Dover Publications, New York, 1972."
```

We shall say that x is **strongly reflexive** in y if the condition `subclass[y, fix[x]]` holds. In the **GOEDEL** program, the (unary) predicate **REFLEXIVE[x]** is defined by the condition

```
In[4]:= subclass[x, cart[fix[x], fix[x]]]
Out[4]= REFLEXIVE[x]
```

In this notebook it is shown that if x is strongly reflexive in y , then the symmetric restriction `intersection[x, cart[y,y]]` is reflexive in the sense of the **GOEDEL** program, but not conversely. The two conditions do agree for the special case that y is the union of the domain and range of x . This is usually called the **field** of x , but to avoid confusion with other meaning of the term field, the following acronym will be introduced for this union:

```
In[5]:= udora[x_] := union[domain[x], range[x]]
```

The restriction `restrict[x, y, z] = intersection[x, cart[y, z]]` of a relation is rewritten by the **GOEDEL** program as a composite:

```
In[6]:= intersection[x, cart[y, z]]
```

```
Out[6]= composite[id[z], x, id[y]]
```

In particular, the symmetric restriction of a class x to the union of its domain and range is the relational part of x , that is, the class of all ordered pairs that belong to x .

```
In[7]:= restrict[x, udora[x], udora[x]]
```

```
Out[7]= composite[Id, x]
```

strong reflexivity => reflexivity of the restriction

Lemma.

```
In[8]:= SubstTest[implies, equal[y, intersection[t, fix[x]]],
  or[equal[0, intersection[y, image[x, intersection[y, complement[fix[x]]]]],
  not[subclass[y, fix[x]]]], t → y] // Reverse
```

```
Out[8]= or[equal[0, intersection[y, image[x, intersection[y, complement[fix[x]]]]],
  not[subclass[y, fix[x]]]] = True
```

```
In[9]:= (% /. {x → x_, y → y_}) /. Equal → SetDelayed
```

Theorem. Strong reflexivity implies reflexivity of the symmetric restriction.

```
In[10]:= or[not[subclass[y, fix[x]]], REFLEXIVE[composite[id[y], x, id[y]]] // AssertTest
```

```
Out[10]= or[not[subclass[y, fix[x]]], REFLEXIVE[composite[id[y], x, id[y]]] = True
```

```
In[11]:= or[not[subclass[y_, fix[x_]]], REFLEXIVE[composite[id[y_], x_, id[y_]]]] := True
```

Counterexample. The converse statement does not hold. The strong assertion `subclass[y, fix[x]]` does not follow from the weaker statement `REFLEXIVE[restrict[x,y,y]]`.

```
In[12]:= (implies[REFLEXIVE[composite[id[y], x, id[y]]], subclass[y, fix[x]]) /.
  {x → EQUIDIFF, y → V}
```

```
Out[12]= False
```

Comment. The statement that x is strongly reflexive on `udora[x]` is equivalent to the statement that the symmetric restriction of x to `udora[x]` is reflexive.

```
In[13]:= subclass[udora[x], fix[x]]
```

```
Out[13]= REFLEXIVE[composite[Id, x]]
```

the class of sets for which the restriction of x is reflexive

In this section some properties of the class of sets on which the symmetric restriction of x is reflexive are derived. The following formula for this class is readily derived: **fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]**. Indeed, one has:

```
In[14]:= member[y, fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]
Out[14]= and[member[y, V], REFLEXIVE[composite[id[y], x, id[y]]]]
```

Lemma. Simplification rule.

```
In[15]:= subclass[intersection[y, image[inverse[x], w]], z] // AssertTest // Reverse
Out[15]= equal[0, intersection[w, image[x, intersection[y, complement[z]]]] ==
          subclass[intersection[y, image[inverse[x], w]], z]

In[16]:= equal[0, intersection[w_, image[x_, intersection[y_, complement[z_]]]] :=
          subclass[intersection[y, image[inverse[x], w]], z]
```

Theorem. The statement that the restriction of x to y is symmetric can be rewritten as follows:

```
In[17]:= SubstTest[subclass, t, cartsq[fix[t]], t → restrict[x, y, y]] // Reverse
Out[17]= and[subclass[intersection[y, image[x, y]], fix[x]],
           subclass[intersection[y, image[inverse[x], y]], fix[x]] ==
           REFLEXIVE[composite[id[y], x, id[y]]]

In[18]:= and[subclass[intersection[y_, image[x_, y_]], fix[x_]],
             subclass[intersection[y_, image[inverse[x], y_]], fix[x_]] :=
             REFLEXIVE[composite[id[y], x, id[y]]]
```

The class of sets y for which each of these inclusions holds is a class of cliques. This explains the following normalization rule:

```
In[19]:= fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]] // Normality // Reverse
Out[19]= intersection[cliques[union[cart[fix[x], V], complement[x]]],
                     cliques[union[cart[fix[x], V], complement[inverse[x]]]] ==
                     fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]

In[20]:= intersection[cliques[union[cart[fix[x_], V], complement[x_]]],
                     cliques[union[cart[fix[x_], V], complement[inverse[x_]]]] :=
                     fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]
```

Theorem. Since classes of cliques are hereditary, the same holds for the class of sets for which the symmetric restriction of x is reflexive.

```
In[23]:= SubstTest[image, inverse[S], intersection[cliques[u], cliques[v]],
  {u -> union[cart[fix[x], V], complement[x]],
   v -> union[cart[fix[x], V], complement[inverse[x]]]} // Reverse

Out[23]= image[inverse[S], fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]] ==
  fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]

In[24]:= image[inverse[S], fix[image[inverse[CART], image[inverse[IMAGE[id[x_]]], RFX]]] :=
  fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]
```

Lemma.

```
In[28]:= SubstTest[Uchains, cliques[t], t -> intersection[x, y] // Reverse

Out[28]= Uchains[intersection[cliques[x], cliques[y]]] = intersection[cliques[x], cliques[y]]

In[29]:= Uchains[intersection[cliques[x_], cliques[y_]]] := intersection[cliques[x], cliques[y]]
```

Theorem. Since classes of cliques are closed under chains of unions, the same holds for the class of sets for which the symmetric restriction of x is reflexive.

```
In[31]:= SubstTest[Uchains, intersection[cliques[u], cliques[v]],
  {u -> union[cart[fix[x], V], complement[x]],
   v -> union[cart[fix[x], V], complement[inverse[x]]]} // Reverse

Out[31]= Uchains[fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]] ==
  fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]

In[32]:= Uchains[fix[image[inverse[CART], image[inverse[IMAGE[id[x_]]], RFX]]] :=
  fix[image[inverse[CART], image[inverse[IMAGE[id[x]]], RFX]]]
```

Suppes' problem 1, part (e) on page 78.

On page 78, problem 1, part (e), the reader is asked to show that a symmetric, transitive relation is reflexive. This statement holds for the predicate **REFLEXIVE** as defined in the **GOEDEL** program.

```
In[33]:= implies[and[SYMMETRIC[x], TRANSITIVE[x]], REFLEXIVE[x]]

Out[33]= True
```

A version of this statement for restrictions of a relation x to a class y also holds, but a corresponding statement fails for the strong version of reflexivity, as shown by a counterexample.

Lemma. (Temporary rewrite rule.)

```
In[34]:= equal[composite[id[y], x, id[y]], composite[id[y], inverse[x], id[y]]] // AssertTest

Out[34]= equal[composite[id[y], x, id[y]], composite[id[y], inverse[x], id[y]]] =
  subclass[composite[id[y], inverse[x], id[y]], x]
```

```
In[35]:= equal[composite[id[y_], x_, id[y_]], composite[id[y_], inverse[x_], id[y_]] :=
  subclass[composite[id[y], inverse[x], id[y]], x]
```

Theorem. If the restriction of a relation is symmetric and transitive, then it is reflexive.

```
In[36]:= Map[implies[#, REFLEXIVE[composite[id[y], x, id[y]]] &,
  SubstTest[and, SYMMETRIC[t], TRANSITIVE[t], t → restrict[x, y, y]] // Reverse
```

```
Out[36]= or[not[subclass[composite[id[y], inverse[x], id[y]], x]],
  not[TRANSITIVE[composite[id[y], x, id[y]]],
  REFLEXIVE[composite[id[y], x, id[y]]] == True
```

```
In[37]:= or[not[subclass[composite[id[y_], inverse[x_], id[y_]], x_]],
  not[TRANSITIVE[composite[id[y_], x_, id[y_]]],
  REFLEXIVE[composite[id[y_], x_, id[y_]]] := True
```

Counterexample. If the restriction of a relation is symmetric and transitive, it need not be reflexive in the stronger sense.

```
In[38]:= (or[not[subclass[composite[id[y], inverse[x], id[y]], x]],
  not[TRANSITIVE[composite[id[y], x, id[y]]],
  subclass[y, fix[x]]]) /. {x → EQUIDIFF, y → V}
```

```
Out[38]= False
```

Corollary. For the case of $y = \text{udora}[x]$, one finds:

```
In[39]:= SubstTest[or, not[subclass[composite[id[y], inverse[x], id[y]], x]],
  not[TRANSITIVE[composite[id[y], x, id[y]]],
  REFLEXIVE[composite[id[y], x, id[y]], y → udora[x]] // Reverse
```

```
Out[39]= or[not[subclass[inverse[x], x]],
  not[TRANSITIVE[composite[Id, x]], REFLEXIVE[composite[Id, x]]] == True
```

```
In[40]:= or[not[subclass[inverse[x_], x_]],
  not[TRANSITIVE[composite[Id, x_]], REFLEXIVE[composite[Id, x_]]] := True
```