

Note on Sunflowers

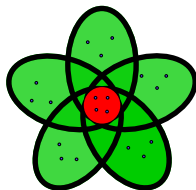
Tolson Bell
Georgia Tech

Suchakree Chueluecha
Lehigh University

Lutz Warnke
Georgia Tech

Based on 2020 REU at Georgia Tech:
<https://arxiv.org/abs/2009.09327>

- Let \mathcal{F} be a k -**uniform** family of subsets of X , i.e., $|S| = k$ and $S \subseteq X$ for all $S \in \mathcal{F}$
- \mathcal{F} is a **sunflower with p petals** if $|\mathcal{F}| = p$ and there exists $Y \subseteq X$ with $Y = S_i \cap S_j$ for all distinct $S_i, S_j \in \mathcal{F}$
- Y is the **core** and $S_i \setminus Y$ are the **petals**
- Note that p disjoint sets forms a sunflower with p petals and empty core.



Sunflower with $k = 7$ and $p = 5$

Applications

Sunflowers have many uses in computer science:

- Fast algorithms for matrix multiplication
- Lower bounds on circuitry
- Cryptography
- Data structure efficiency
- Pseudorandomness
- Random approximations

Research Question

What is the smallest $r = r(p, k)$ such that every k -uniform family with at least r^k sets must contain a sunflower with p petals?

Erdős–Rado (1960)

- (a) $r = pk$ is **sufficient** to guarantee a sunflower:
every family with more than $(pk)^k > k!(p-1)^k$ sets contains a sunflower
- (b) $r > p-1$ is **necessary** to guarantee a sunflower:
there is a family of $(p-1)^k$ sets without a sunflower

- Erdős conjectured $r = r(p)$ is sufficient (no k dependency), offered \$1000 reward
- Until 2018, best known upper bound on r was still $k^{1-o(1)}$ with respect to k

"[The sunflower problem] has fascinated me greatly – I really do not see why this question is so difficult."

–Paul Erdős (1981)

- Erdős conjectured $r = r(p)$ is sufficient (no k dependency)
- Until 2018, best known upper bound on r was still $k^{1-o(1)}$ with respect to k

Alweiss–Lovett–Wu–Zhang (Breakthrough Aug 2019)

$r = p^3(\log k)^{1+o(1)}$ is sufficient to guarantee a sunflower

New papers built off their breakthrough ideas:

- Sep 2019: *Rao* used Shannon's coding theorem for a cleaner proof and slightly better bound
- Oct 2019: *Frankston–Kahn–Narayanan–Park* improved a key lemma of ALWZ, enabling them to prove a conjecture of Talagrand regarding thresholds functions
- Jan 2020: *Rao* improved to $r = O(p \log(pk))$ by incorporating ideas from FKNP
- July 2020: *Tao* matched Rao's bound with shorter proof using Shannon entropy

Rao (Jan 2020)

$r = O(p \log(pk))$ is sufficient to guarantee a sunflower

Bell–Chueluecha–Warnke (September 2020)

$r = O(p \log k)$ is sufficient to guarantee a sunflower

Further REU 2020 results:

- *Rao/Tao methods not needed for this result:*
2019 Frankston–Kahn–Narayanan–Park result suffices with our proof variant
- *Main Technical Lemma is asymptotically sharp:*
Bound cannot be improved further without change of proof strategy

Strategy: Reduction to r -spread Families

- *Key Definition:* \mathcal{F} is **r -spread** if $|\mathcal{F}| \geq r^k$ and for every nonempty $S \subseteq X$ the number of sets in \mathcal{F} which contain S is at most $r^{k-|S|}$

The Inductive Reduction

If every r -spread family contains p disjoint sets, then r^k sets guarantees a sunflower.

Proof. Induction on k .

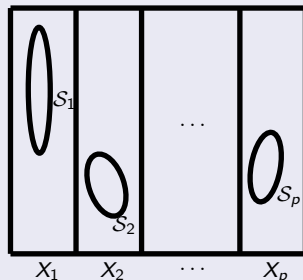
Question: How to *find* p disjoint sets in an r -spread family?

- We now review the common proof framework of previous work.

Question: How to find p disjoint sets?

The Probabilistic Method

- 1 Consider a random partition of X to X_1, X_2, \dots, X_p ($x \in X$ goes in random X_i)
- 2 Use **probabilistic method**
 - Show $\mathbb{P}(\nexists S_i \in \mathcal{F} \text{ such that } S_i \subseteq X_i) < \frac{1}{p}$
 - *Union bound:*
 $\mathbb{P}(\exists i \text{ where } X_i \text{ has no } S_i) < p \cdot \frac{1}{p} = 1$
 - There is partition where each X_i has S_i
 - Then S_1, \dots, S_p are disjoint sets in \mathcal{F}



Main Technical Lemma (Rao 2020)

Let X_p be set where $\forall x \in X, x \in X_p$ w.p. $\frac{1}{p}$ independently. $\exists C > 1$ s.t. for $r \geq Cp \log(pk)$, $\mathbb{P}(\text{There does not exist } S_i \in \mathcal{F} \text{ such that } S_i \subseteq X_p) < \frac{1}{p}$

Main Technical Lemma (Rao 2020)

Let X_a be set where $\forall x \in X, x \in X_a$ w.p. $\frac{1}{a}$ independently. $\exists C > 1$ s.t. for $r \geq C a \log(bk)$, $\mathbb{P}(\text{There does not exist } S_i \in \mathcal{F} \text{ such that } S_i \subseteq X_a) < \frac{1}{b}$

Bell–Chueluecha–Warnke (September 2020)

$r = O(p \log k)$ is sufficient to guarantee a sunflower

Proof Sketch (improve union bound via linearity of expectation):

- Partition X_1, \dots, X_{2p} instead of X_1, \dots, X_p .
- To get p disjoint sets, half of our sets need to contain a set in \mathcal{F}
- *Linearity of expectation*: if each X_i has less than half chance of failure, there is some partition where at least half succeed
- Apply main lemma with $a = 2p$, $b = 2$.
- $r = 2Cp \log(2k) = O(p \log k)$ suffices!

\mathcal{F} , a k -uniform family of subsets of X , is a **sunflower with p petals** if $|\mathcal{F}| = p$ and there exists $Y \subseteq X$ with $Y = S_i \cap S_j$ for all distinct $S_i, S_j \in \mathcal{F}$.

Research Question

What is the smallest $r = r(p, k)$ such that every k -uniform family with at least r^k sets must contain a sunflower with p petals?

- Erdős–Rado (1960): $r = pk$ is sufficient and $r > p - 1$ is necessary
- Erdős (1981): Conjectured $r = r(p)$ sufficient
- Alweiss–Lovett–Wu–Zhang (2019): Breakthrough that $r = p^3(\log k)^{1+o(1)}$ suffices
- Rao (2020): By Shannon's Coding Theorem, $r = O(p \log(pk))$ suffices

Bell–Chueluecha–Warnke (2020)

- $r = O(p \log k)$ suffices by minor variant of existing probabilistic arguments
- This bound cannot be improved without change of strategy

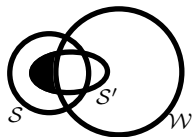
- Alweiss–Lovett–Wu–Zhang (2020). *Improved bounds for the sunflower lemma*. Proceedings of STOC 2020. Extended preprint at arXiv:1908.08483
- Erdős (1981). *On the combinatorial problems which I would most like to see solved*. Combinatorica.
- Erdős–Rado (1960). *Intersection theorems for systems of sets*. Journal of the London Mathematical Society.
- Frankston–Kahn–Narayanan–Park (2019). *Thresholds versus fractional expectation-thresholds*. Preprint at arXiv:1910.13433.
- Rao (2020). *Coding for sunflowers*. Discrete Analysis. Preprint at arXiv:1909.04774
- Tao (2020). *The sunflower lemma via Shannon entropy*. terrytao.wordpress.com.
- Bell–Chueluecha–Warnke (2020). *Note on Sunflowers*. Discrete Mathematics. Preprint at arXiv:2009.09327

Main Technical Lemma

$\mathbb{P}(\text{There \underline{does not} exist } S_i \in \mathcal{F} \text{ such that } S_i \subseteq X_a) < \frac{1}{b}$

Proof. Partition X_i to V_1, V_2 with equal size, so $|V_1| = |V_2| = |X|/(2a)$.

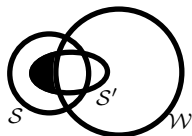
- *Key Definition:* Given $S \in \mathcal{F}$ and $W \subseteq X$, (S, W) is **m-good** if there exists $S' \in \mathcal{F}$ such that $S' \subseteq W \cup S$ and $|S' \setminus W| \leq m$



Iteration: $\mathbb{P}(\text{Less than half of sets in } \mathcal{F} \text{ are } m\text{-good with respect to } V_1) \leq \frac{1}{2b}$

Final Step: $\mathbb{P}(V_1 \cup V_2 \text{ does not contain a set in } \mathcal{F} \mid \text{successful iteration}) < \frac{1}{2b}$

- *Key Definition:* Given $S \in \mathcal{F}$ and $W \subseteq X$, (S, W) is **m -good** if there exists $S' \in \mathcal{F}$ such that $S' \subseteq W \cup S$ and $|S' \setminus W| \leq m$



Iteration: $\mathbb{P}(\text{Less than half of sets in } \mathcal{F} \text{ are } m\text{-good with respect to } V_1) \leq \frac{1}{2b}$

- Partition V_1 to W_1, W_2, \dots, W_x with equal size
- Iteratively replace each good $(S, \bigcup_{1 \leq i \leq j} W_i)$ pair with the guaranteed S'
- Bound the number of bad pairs by a key counting lemma & Markov's inequality
- Moving from S to S' reduces the set sizes at each step as $\bigcup_{1 \leq i \leq j} W_i$ expands

Final Step: $\mathbb{P}(V_1 \cup V_2 \text{ does not contain a set in } \mathcal{F} \mid \text{successful iteration}) < \frac{1}{2b}$

- Construct an m -uniform \mathcal{F}' from sets in \mathcal{F} which are m -good with respect to V_1
- Apply Janson's Inequality with V_2 and \mathcal{F}' to bound $\mathbb{P}(\exists S \in \mathcal{F}' \text{ s.t. } S \subseteq V_2)$