

## CS 1050 Homework 9 Solutions

1. Consider  $c = 3$ , then we have,

$$c \cdot g(n) = 3 \cdot 2^n > 2^{n+1} \text{ for all } n \geq 1.$$

$$\Leftrightarrow f(n) < c \cdot g(n) \text{ for all } n \geq 1.$$

$$\Leftrightarrow f(n) = O(g(n))$$

2.a We have,

$$\frac{f(n)}{g(n)} = \frac{(n+1)^2}{n^2}$$

$$= \frac{n^2 + 2n + 1}{n^2} = 1 + \frac{2}{n} + \frac{1}{n^2}$$

$$\Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} 1 + \frac{2}{n} + \frac{1}{n^2}$$

$$\Leftrightarrow \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \lim_{n \rightarrow \infty} 1 + \lim_{n \rightarrow \infty} \frac{2}{n} + \lim_{n \rightarrow \infty} \frac{1}{n^2}$$

Now, since  $\lim_{n \rightarrow \infty} \frac{2}{n} = 0$  and  $\lim_{n \rightarrow \infty} \frac{1}{n^2} = 0$ , we get

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 1$$

$$\Leftrightarrow f(n) = O(g(n))$$

b. Consider  $c = 4$ , and we have,

$$c \cdot g(n) = 4n^2 = n^2 + 3n^2$$

Since we know that  $1 \leq n \leq n^2$  for all  $n \geq 1$ , we have,

$$c \cdot g(n) \geq n^2 + 2n + 1 = (n+1)^2 \text{ for all } n \geq 1$$

$$\Leftrightarrow f(n) \leq c \cdot g(n) \text{ for all } n \geq 1$$

$$\Leftrightarrow f(n) = O(g(n)).$$

3. We will prove this without using the limit. Clearly,

$$n \geq (200)^2 + 50 \text{ for all } n \geq (200)^2 + 50.$$

$$\begin{aligned}
&\Leftrightarrow n^3 \geq (200n)^2 + 50n^2 \text{ for all } n \geq (200)^2 + 50. \\
&\Leftrightarrow n^3 \geq (200n)^2 + 50 \text{ for all } n \geq (200)^2 + 50 \text{ (using } 50n^2 \geq 50). \\
&\Leftrightarrow g(n) \geq f(n) \text{ for all } n \geq (200)^2 + 50. \\
&\Leftrightarrow f(n) = O(g(n)).
\end{aligned}$$

4. We have,

$$\begin{aligned}
&g(n) - 3f(n) = n^2 + 5n - 6 = (n - 1)(n + 6) \geq 0 \text{ for all } n \geq 1 \\
&\Leftrightarrow g(n) \geq 3f(n) \text{ for all } n \geq 1 \\
&\Leftrightarrow f(n) \leq \frac{1}{3}g(n) \text{ for all } n \geq 1 \\
&\Leftrightarrow f(n) = O(g(n))
\end{aligned}$$

5. **Proof:** Consider  $N = 2A^2$ . Now for any  $n \geq N = 2A^2$ ,

$$n! = A^2!(A^2 + 1) \dots 2A^2 \dots n$$

Now  $A^2! \geq 1$  (since  $A$  is a positive integer) and each term on the RHS after  $A^2!$  is greater than  $A^2$  (there  $n - A^2$  such terms). Therefore,

$$n! > (A^2)^{n-A^2} = A^{2(n-A^2)}$$

Now since  $n \geq 2A^2$ , we have that  $n - \frac{n}{2} \geq A^2$  and so  $n - A^2 \geq \frac{n}{2}$ . Therefore,

$$\begin{aligned}
&n! > A^{2(\frac{n}{2})} \\
&\Leftrightarrow n! > A^n
\end{aligned}$$

So,  $N = 2A^2$  suffices.

6a. We need to prove that  $f$  is not  $O(g)$ . We prove it by contradiction. Suppose that  $f = O(g)$ . That means that there exists constants  $c$  and  $n_0$  such that,

$$c \cdot g(n) \geq f(n) \text{ for all } n \geq n_0$$

Now pick  $n_1$  to be any even number greater than  $n_0$  and  $c$ , i.e.  $n_1 > \max\{n_0, c\}$  and  $n_1$  is even. Clearly such an even number can be picked since  $c$  and  $n_0$  are constants. We have,

$$f(n_1) = n_1^2 > c n_1 = c \cdot g(n_1)$$

which is a contradiction to the fact that  $c \cdot g(n) \geq f(n)$  for all  $n \geq n_0$  (since

$n_1 > n_0$  and  $n_1 > c$ )

Therefore  $f(n)$  cannot be  $O(g(n))$ .